

# COMPILER OPTIMIZATION FOR (OPENMP) ACCELERATOR OFFLOADING

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EuroLLVM – April 8, 2019 – Brussels, Belgium

Johannes Doerfert and Hal Finkel

Leadership Computing Facility  
Argonne National Laboratory  
<https://www.alcf.anl.gov/>



## ACKNOWLEDGMENT

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This research was supported by the Exascale Computing Project (17-SC-20-SC), a collaborative effort of two U.S. Department of Energy organizations (Office of Science and the National Nuclear Security Administration) responsible for the planning and preparation of a capable exascale ecosystem, including software, applications, hardware, advanced system engineering, and early testbed platforms, in support of the nation's exascale computing imperative.



## COMPILER OPTIMIZATION

Original Program

```
int y = 7;  
  
for (i = 0; i < N; i++) {  
    f(y, i);  
}  
g(y);
```

After Optimizations

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# COMPILER OPTIMIZATION FOR PARALLELISM

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## CURRENT COMPILER OPTIMIZATION FOR PARALLELISM

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None<sup>\*†</sup>

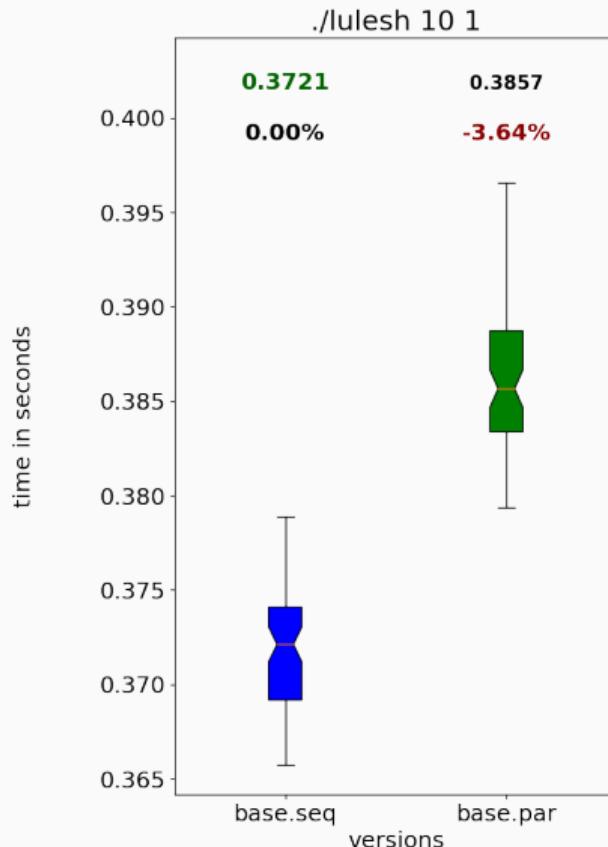
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\*At least for LLVM/Clang up to 8.0

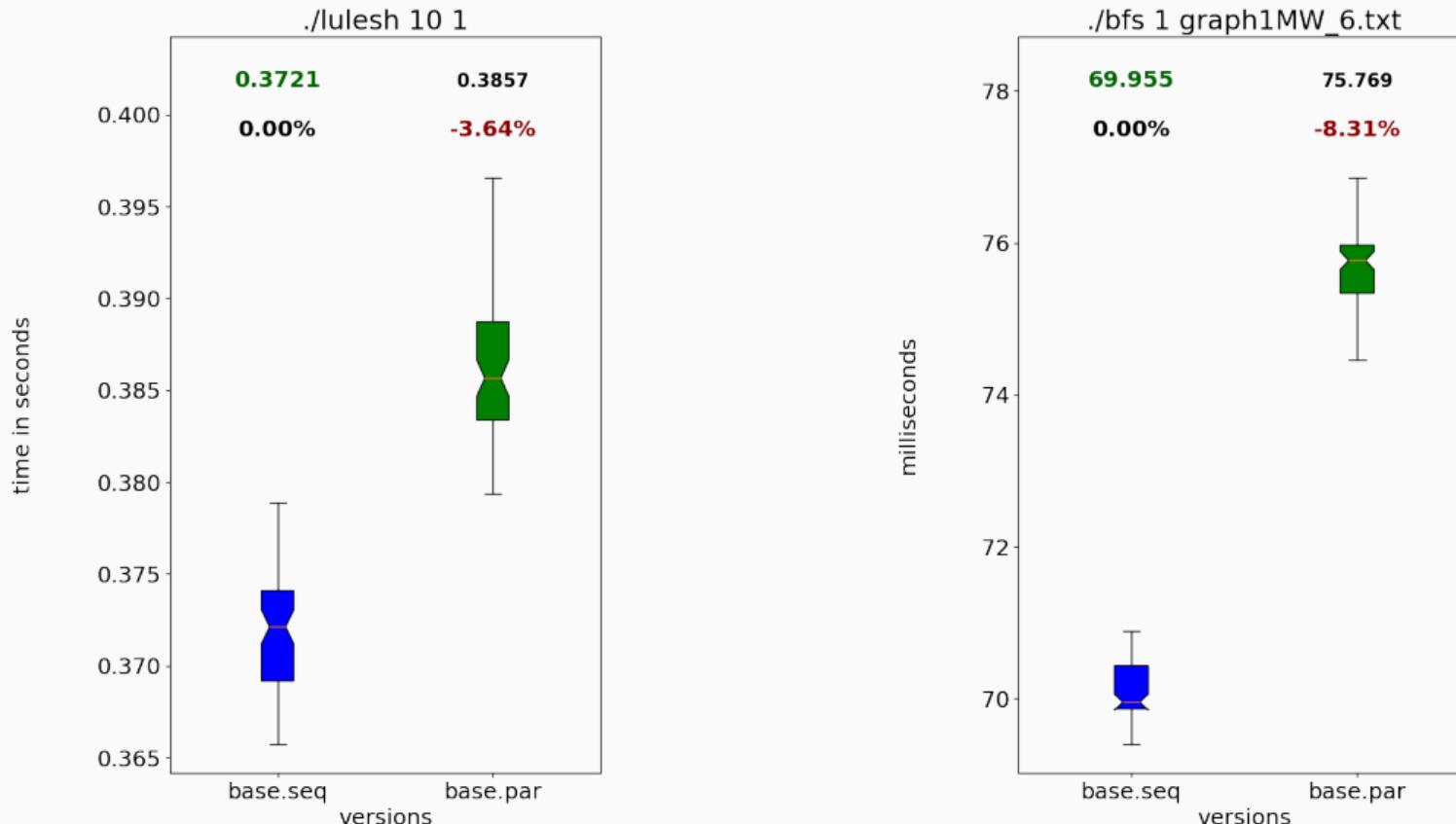
†And not considering smart **runtime libraries**!

Why is this important?

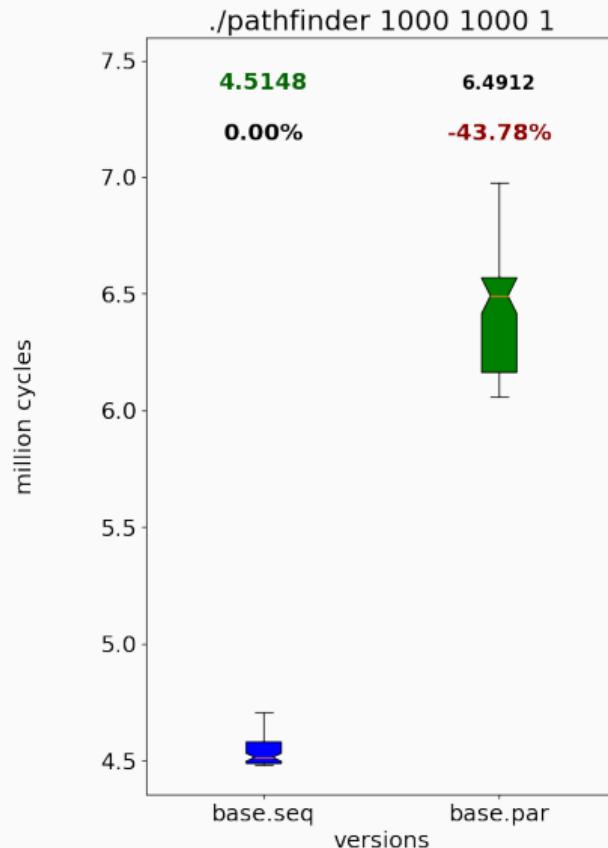
# PERFORMANCE IMPLICATIONS



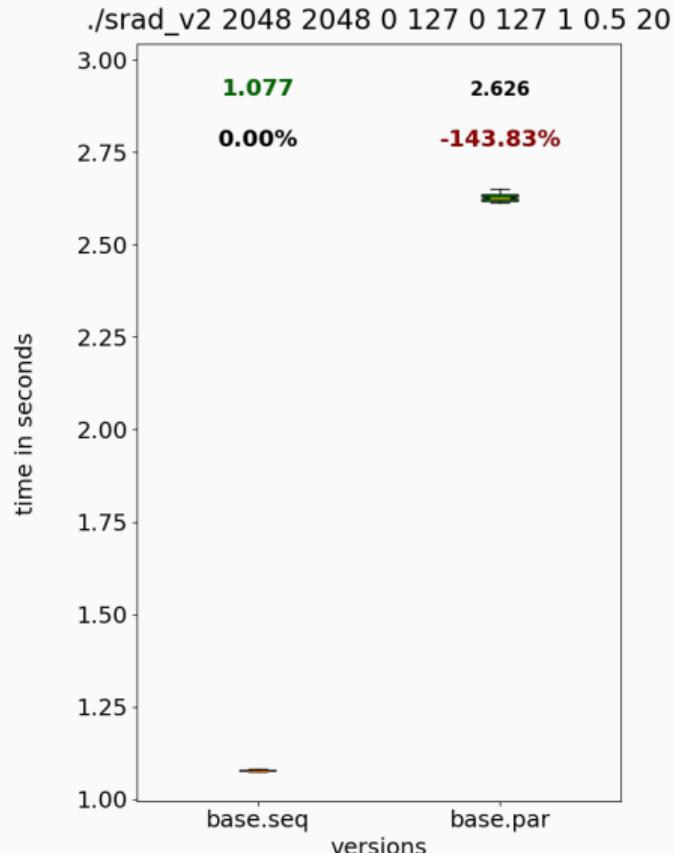
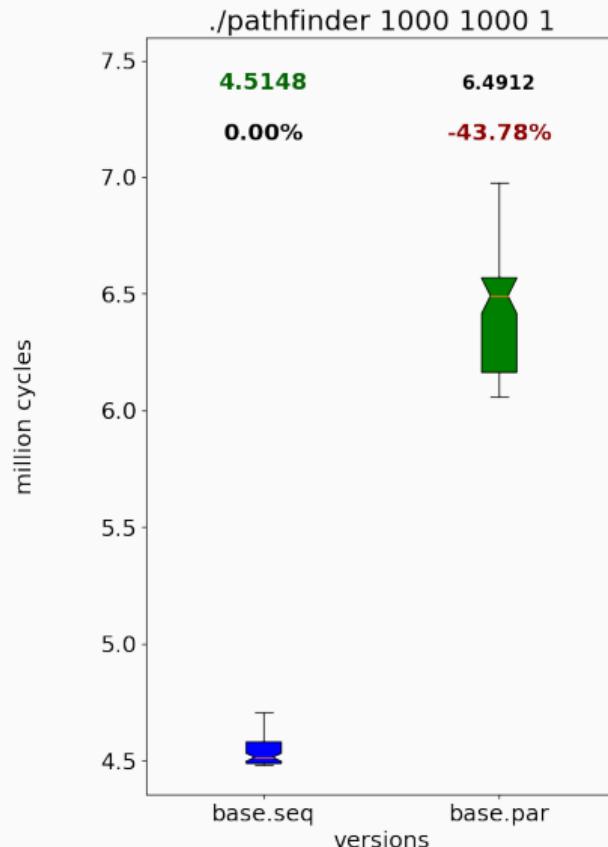
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## OPTIMIZATION CATEGORIES

Optimizations for *sequential* aspects

Optimizations for *parallel* aspects

---

a  
b



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- May reuse *existing* transformations (patches up for review!)

### Optimizations for *parallel* aspects

- New *explicit parallelism-aware* transformations

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    #pragma omp target teams  
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    work();                                // <- Hotspot  
    #pragma omp target teams  
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N teams, with M threads each, all executing `work` concurrently.



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OpenMP

Clang

LLVM-IR

LLVM

Assembly

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Assembly

Code



# OPENMP OFFLOAD — OVERVIEW

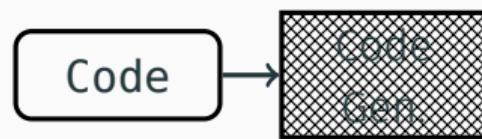
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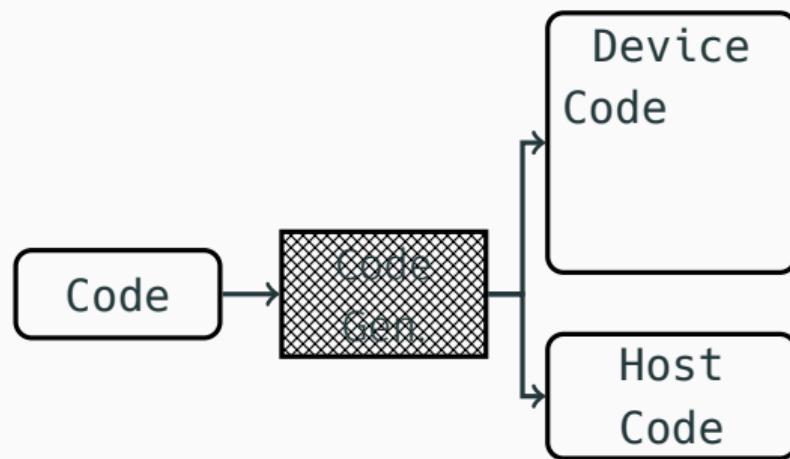
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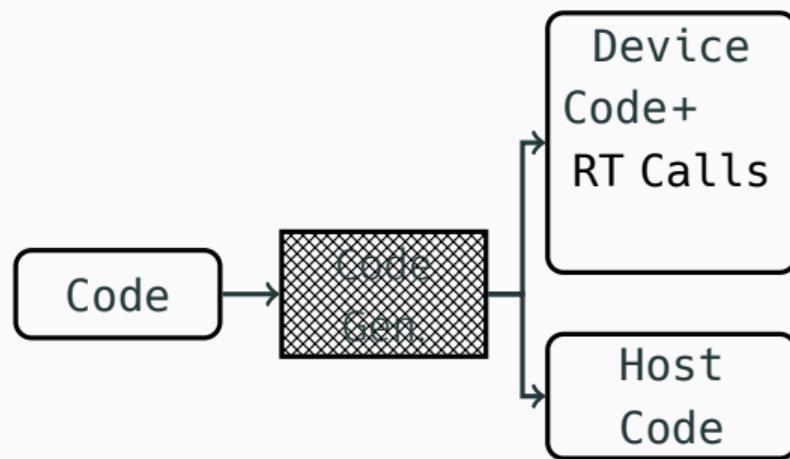
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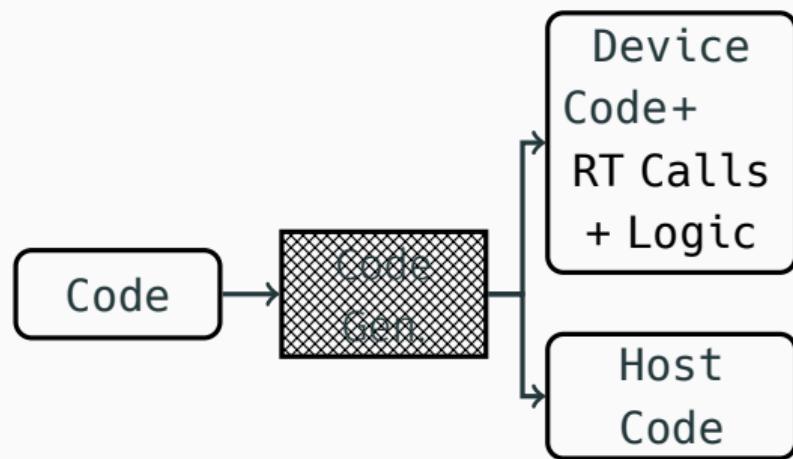
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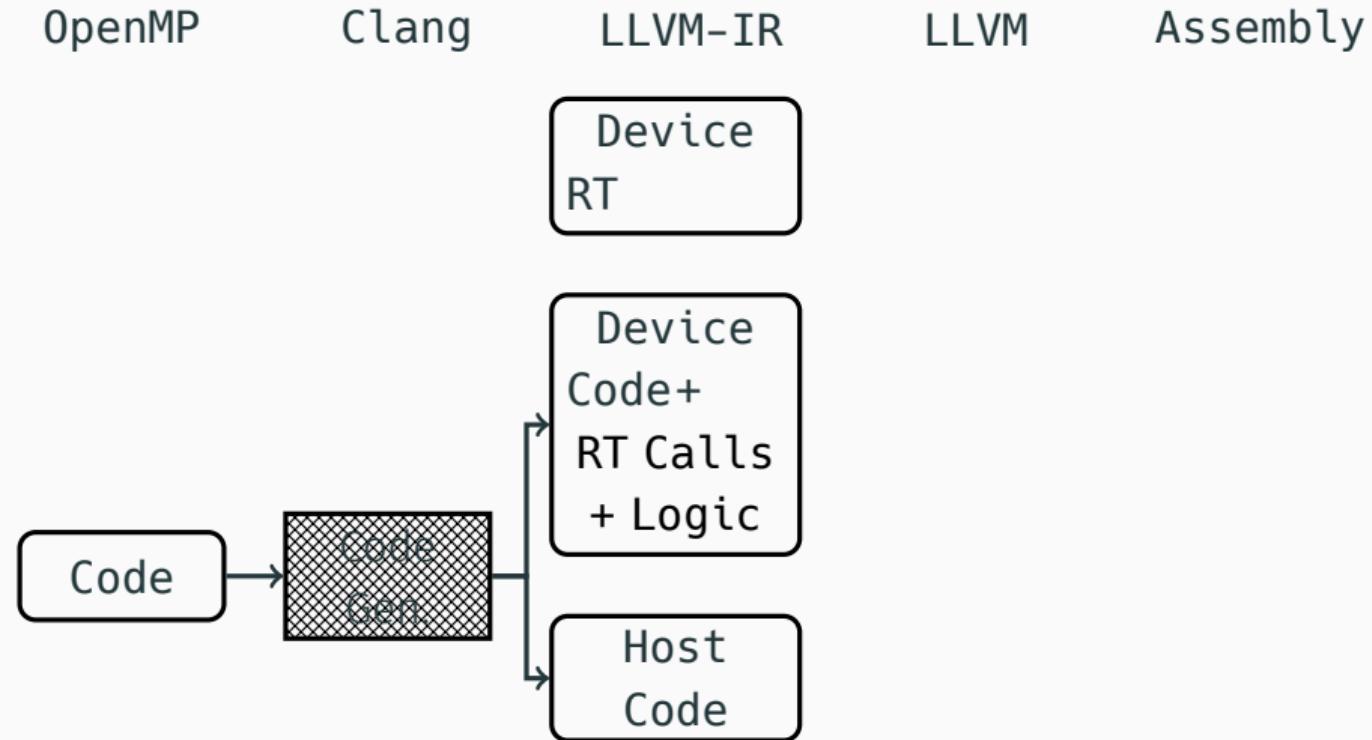


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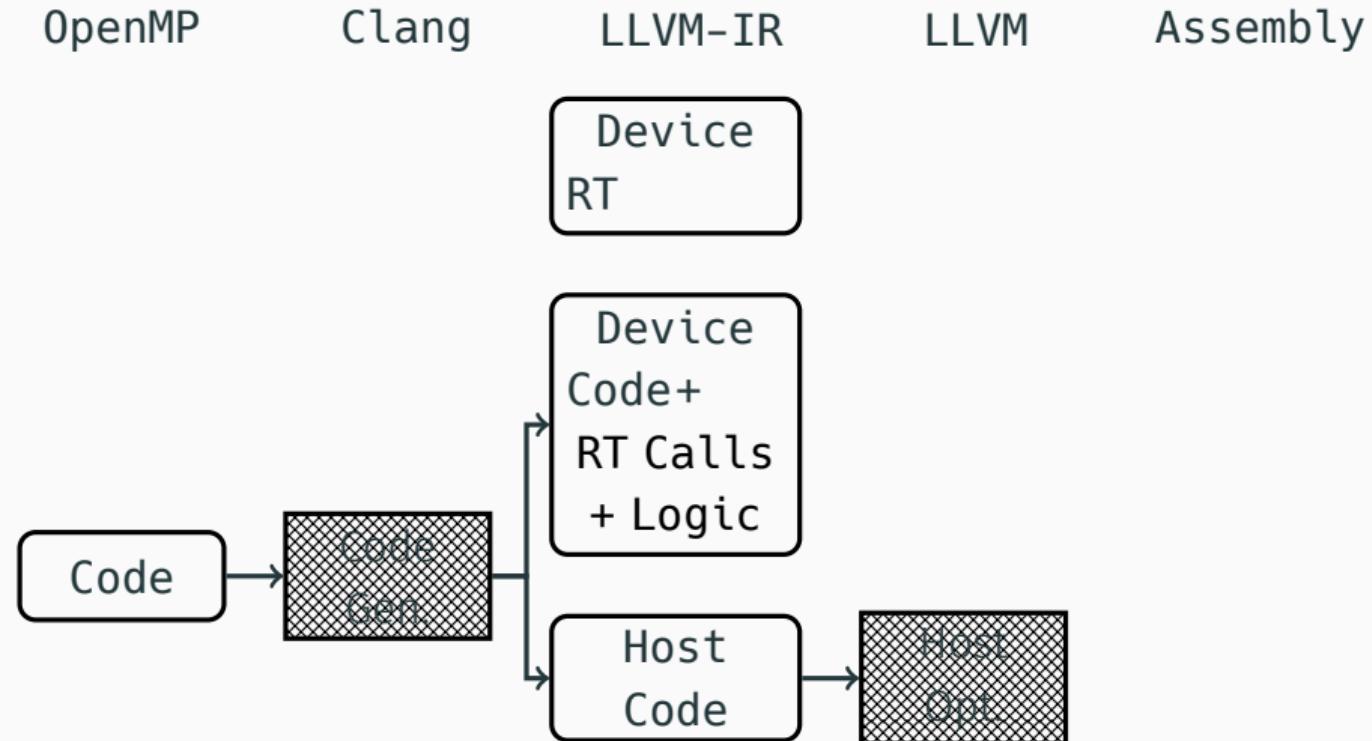
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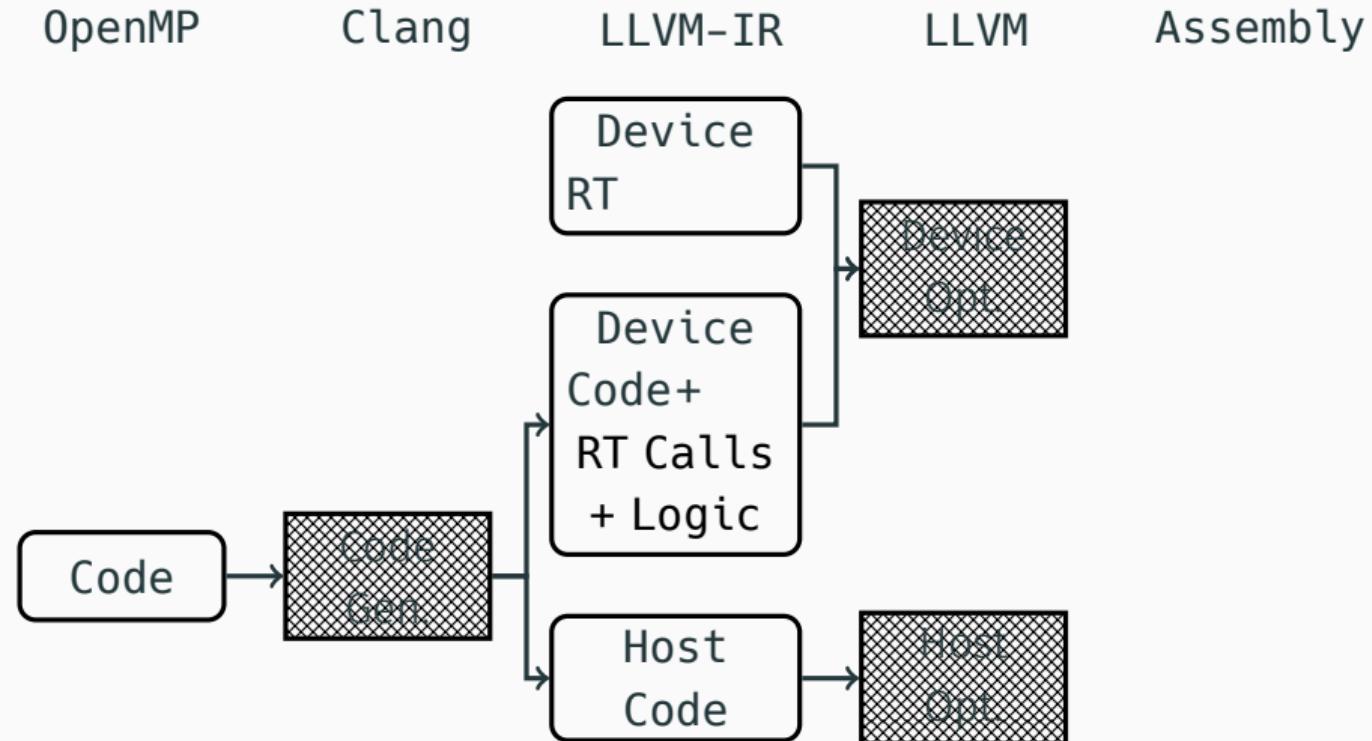
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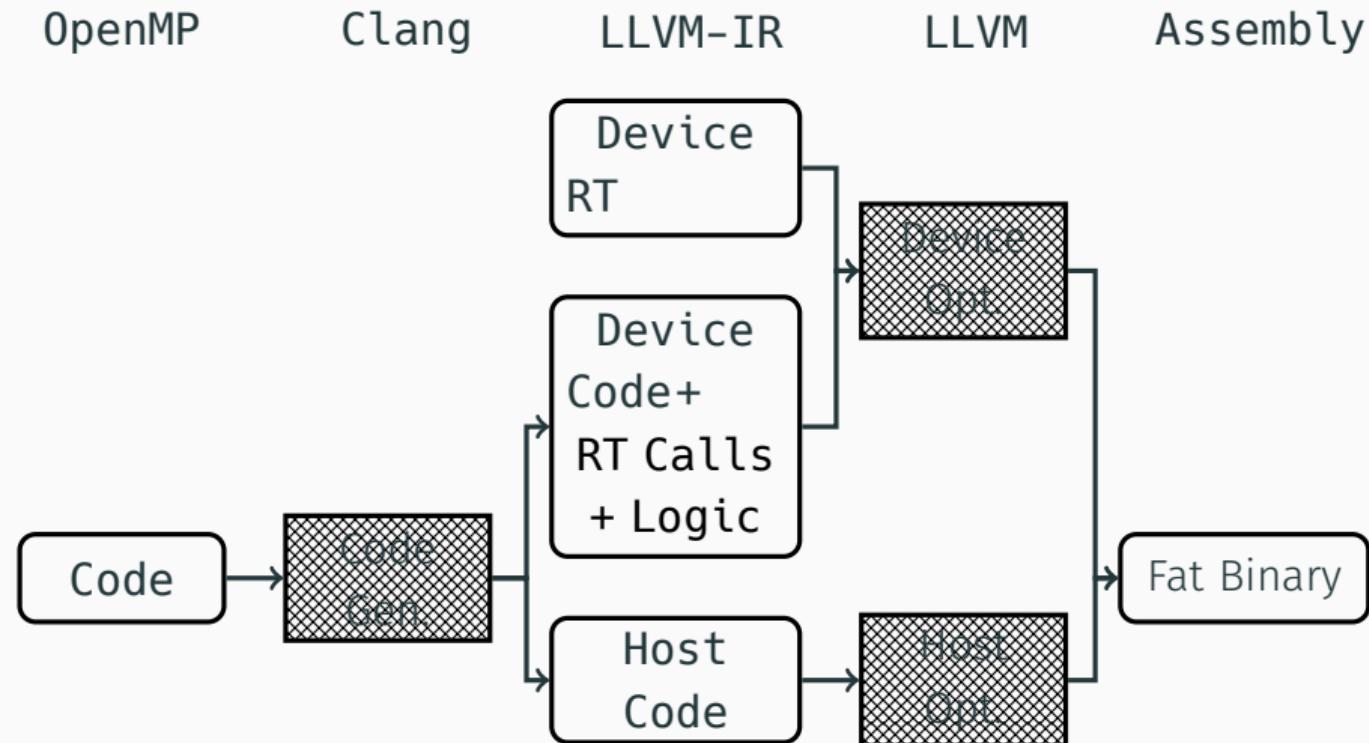
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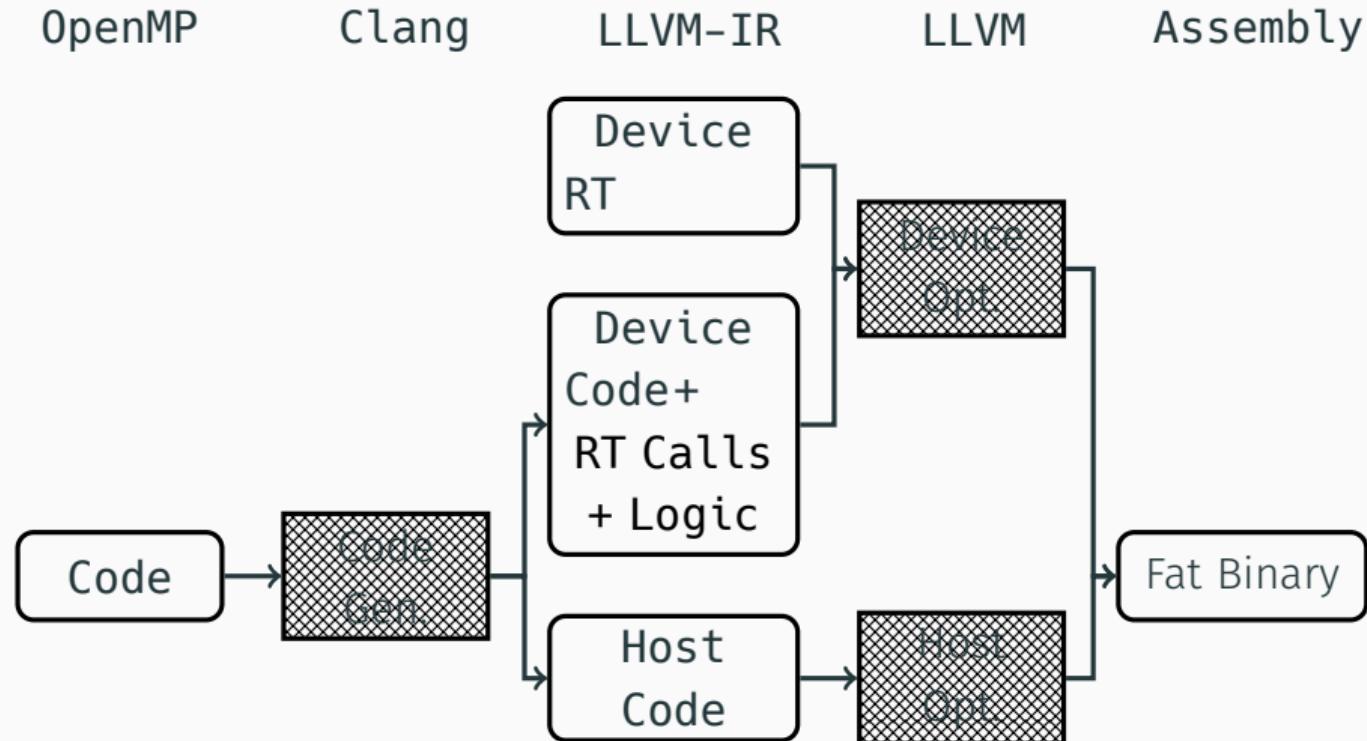
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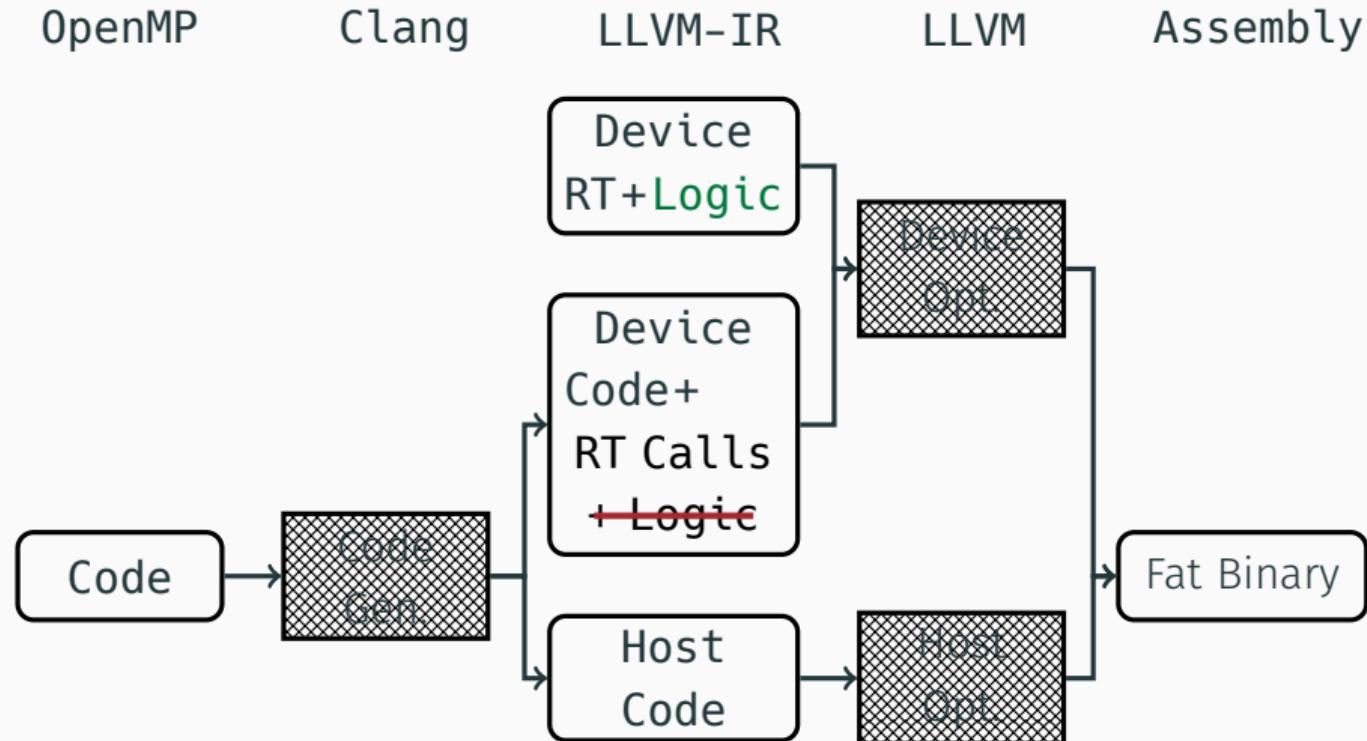
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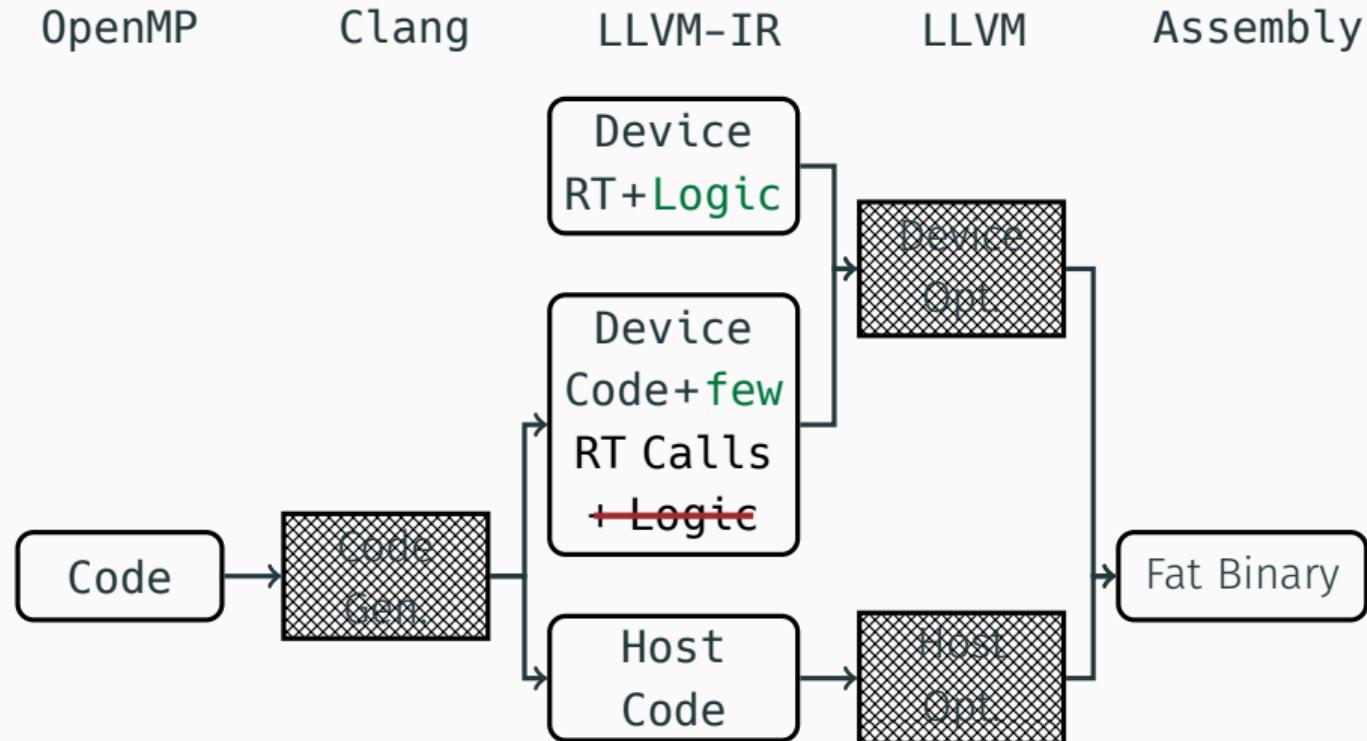
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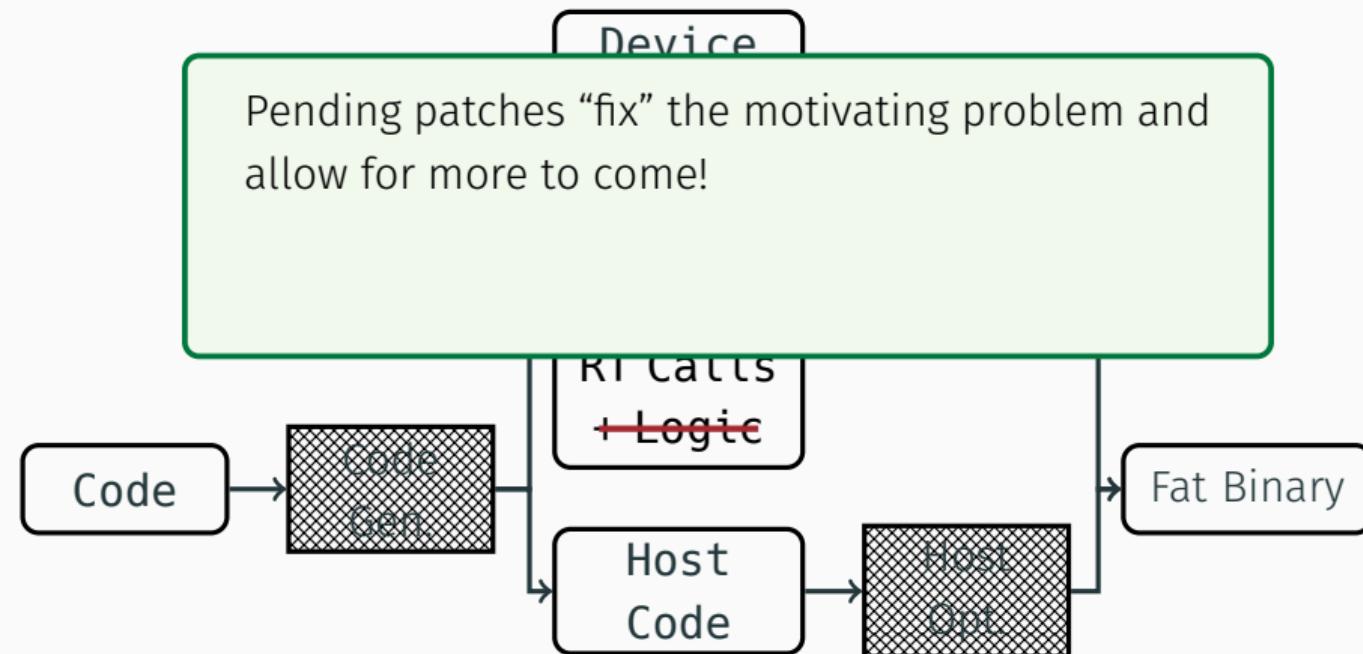


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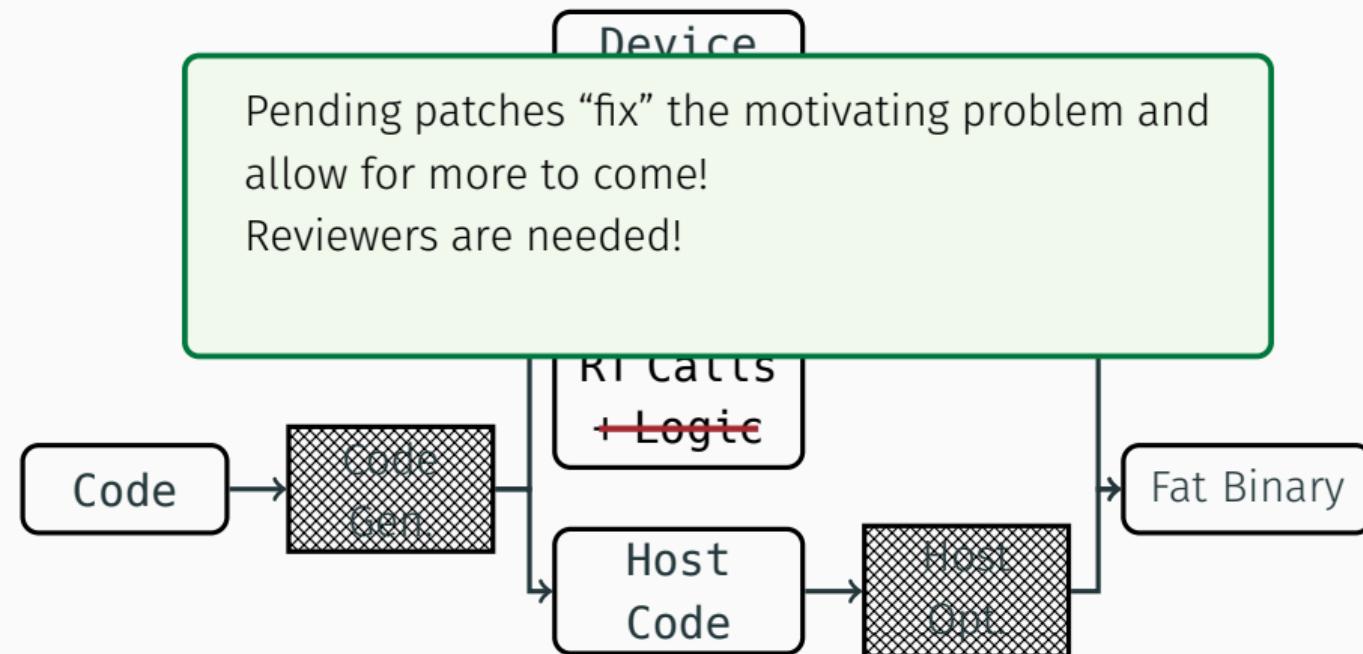
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OpenMP      Clang      LLVM-IR      LLVM      Assembly



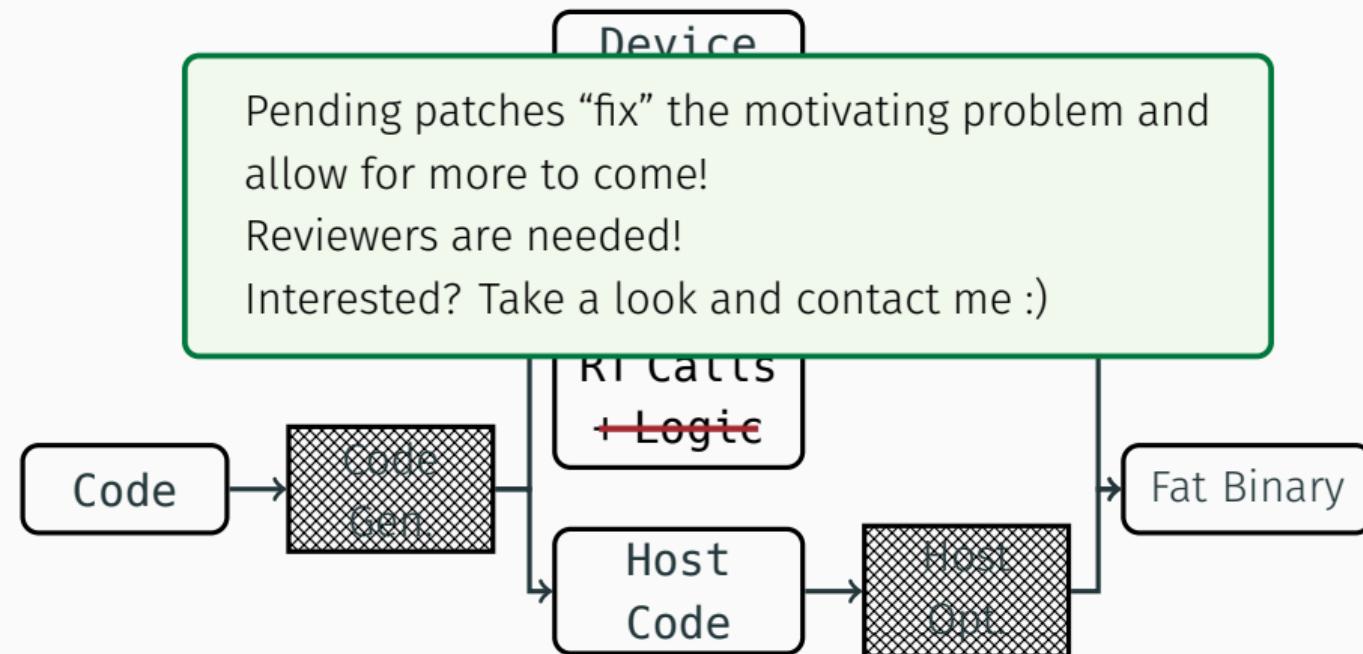
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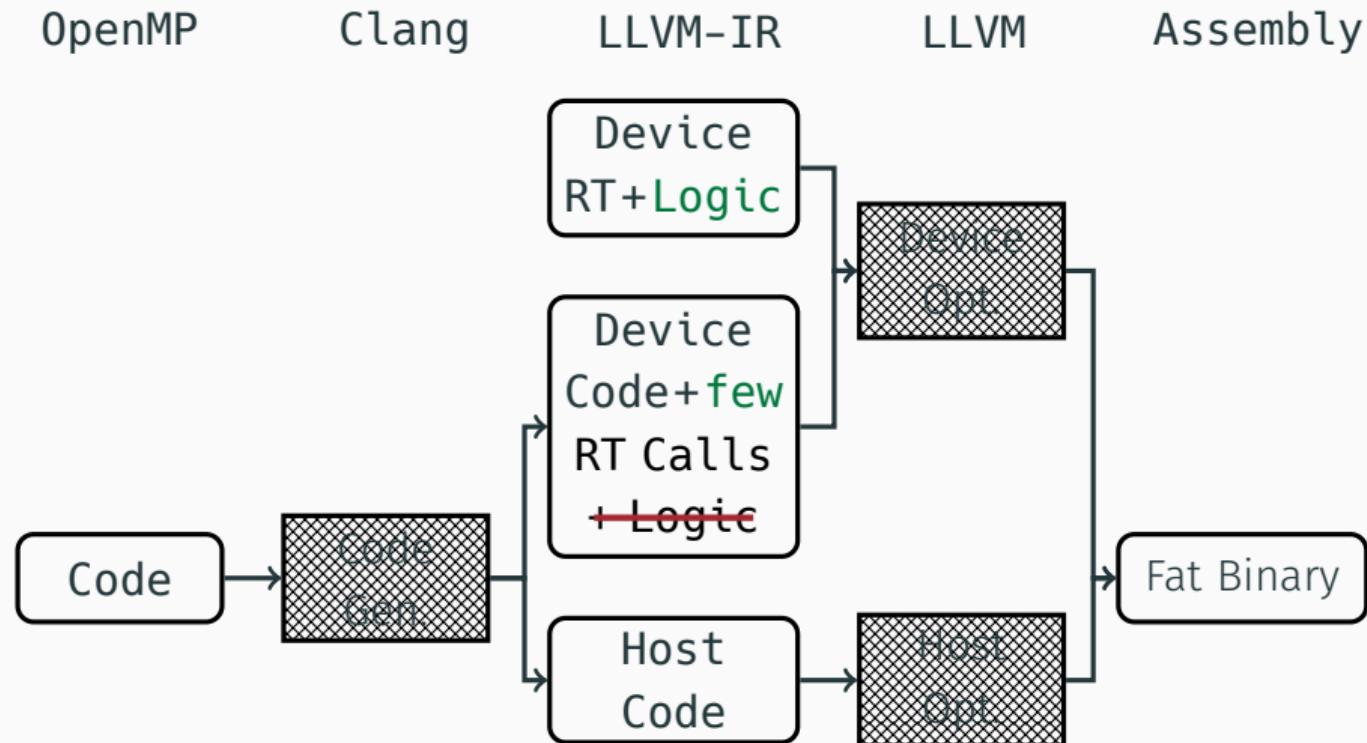


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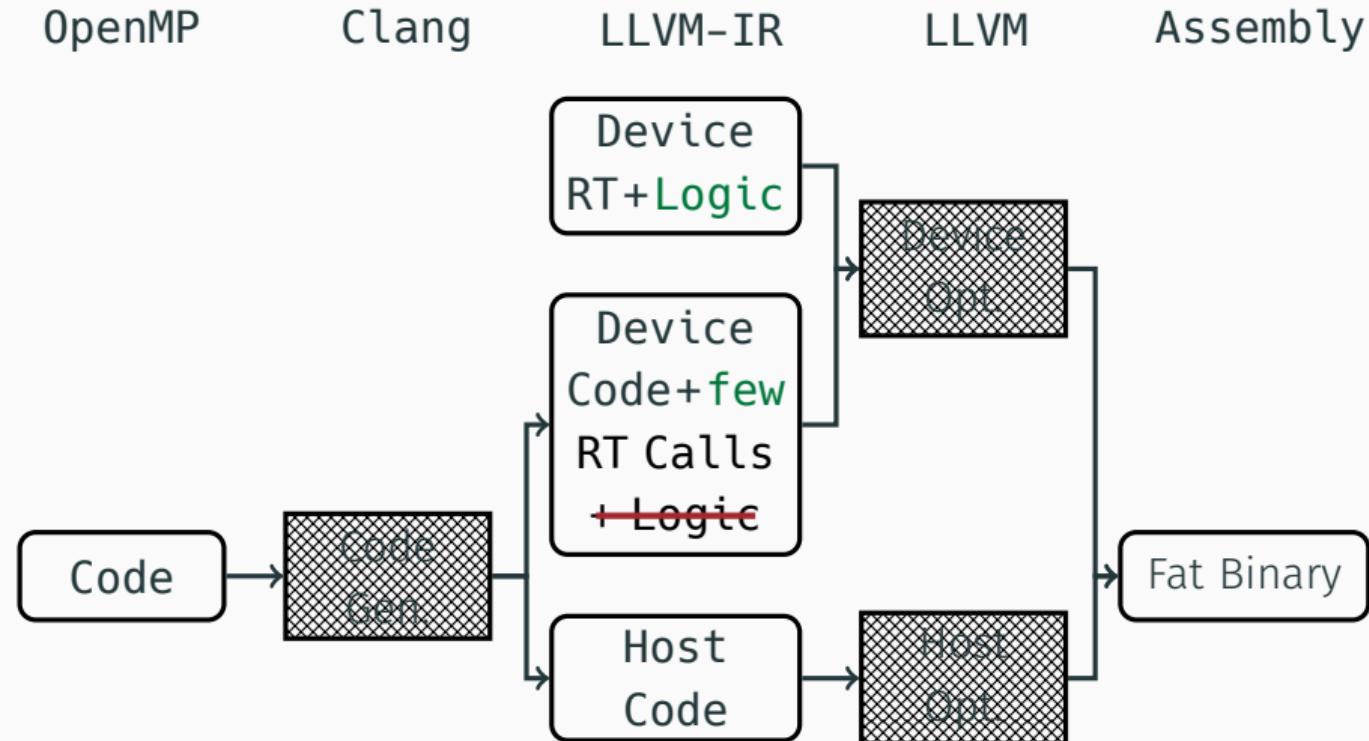
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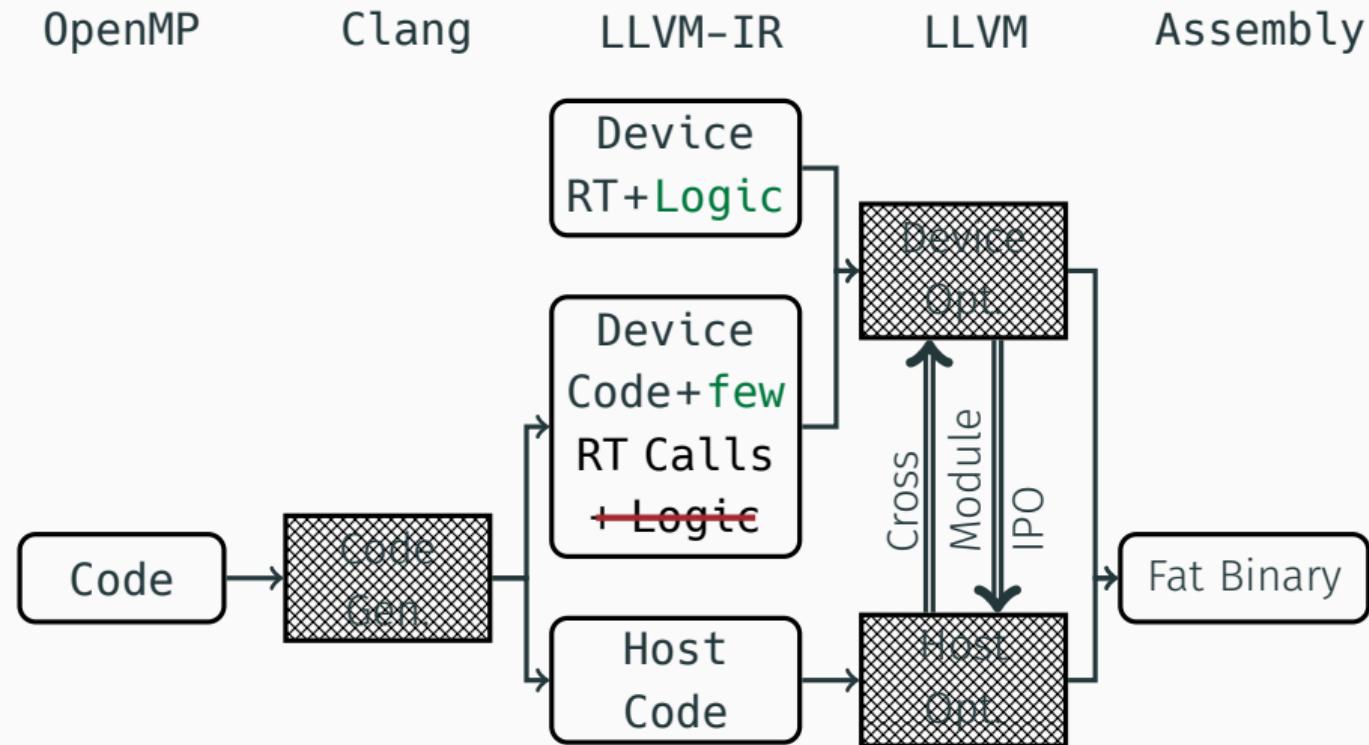
# OPENMP OFFLOAD — OVERVIEW & DIRECTIONS



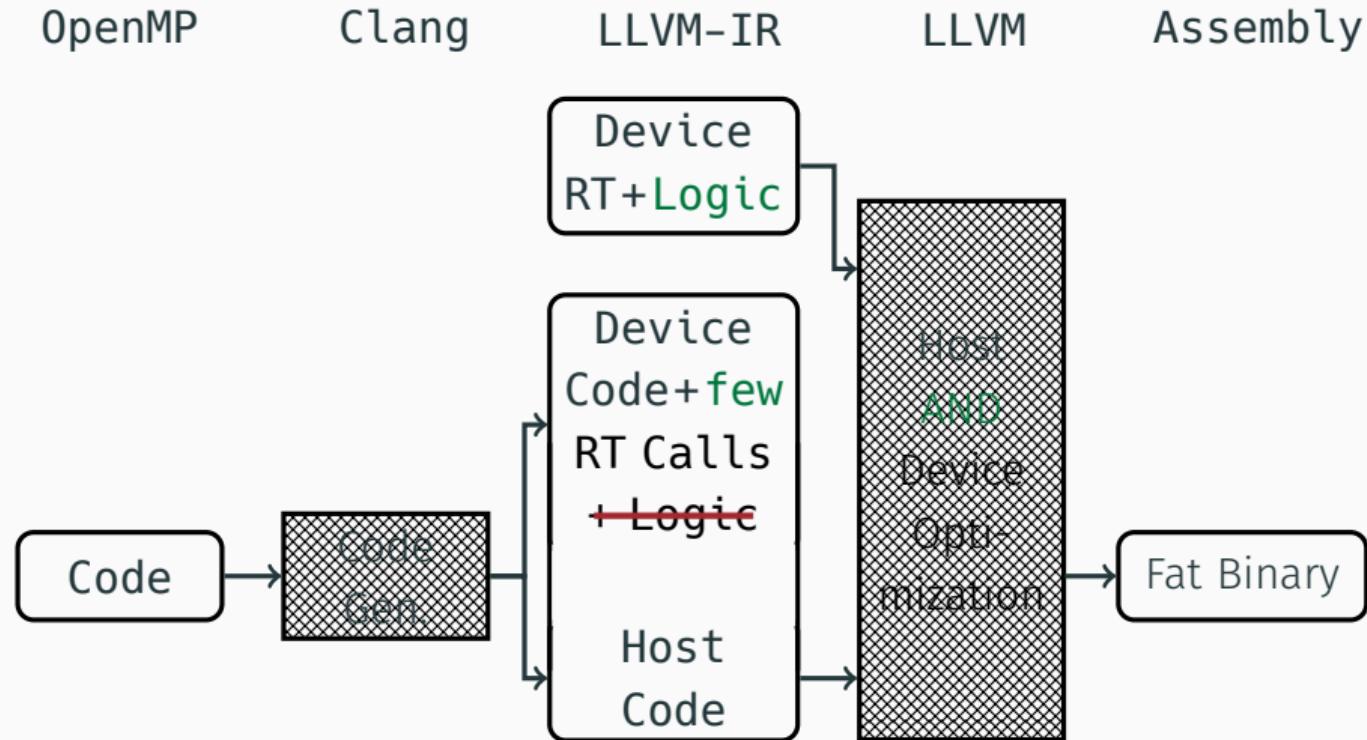
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- Device RT interface & implementation are separated:
  - ◊ simplifies generated LLVM-IR
  - ◊ most LLVM & Clang parts become target agnostic

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- use inter-procedural reasoning to place minimal guards/synchronization
- if legal, switch all boolean **UseSPMDMode** flags to **true**
- currently, no (unknown) global side-effects allowed outside parallel regions.



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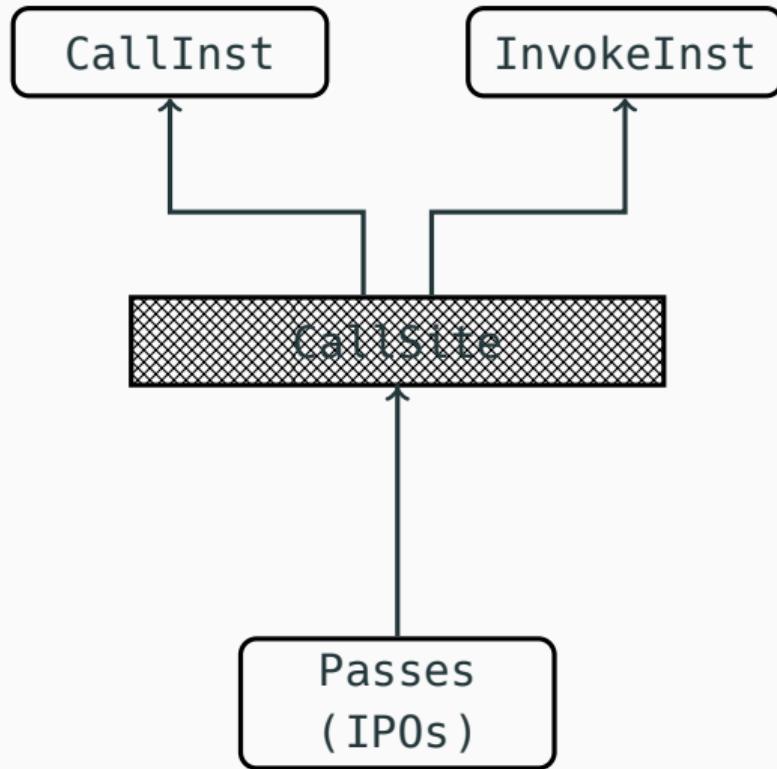


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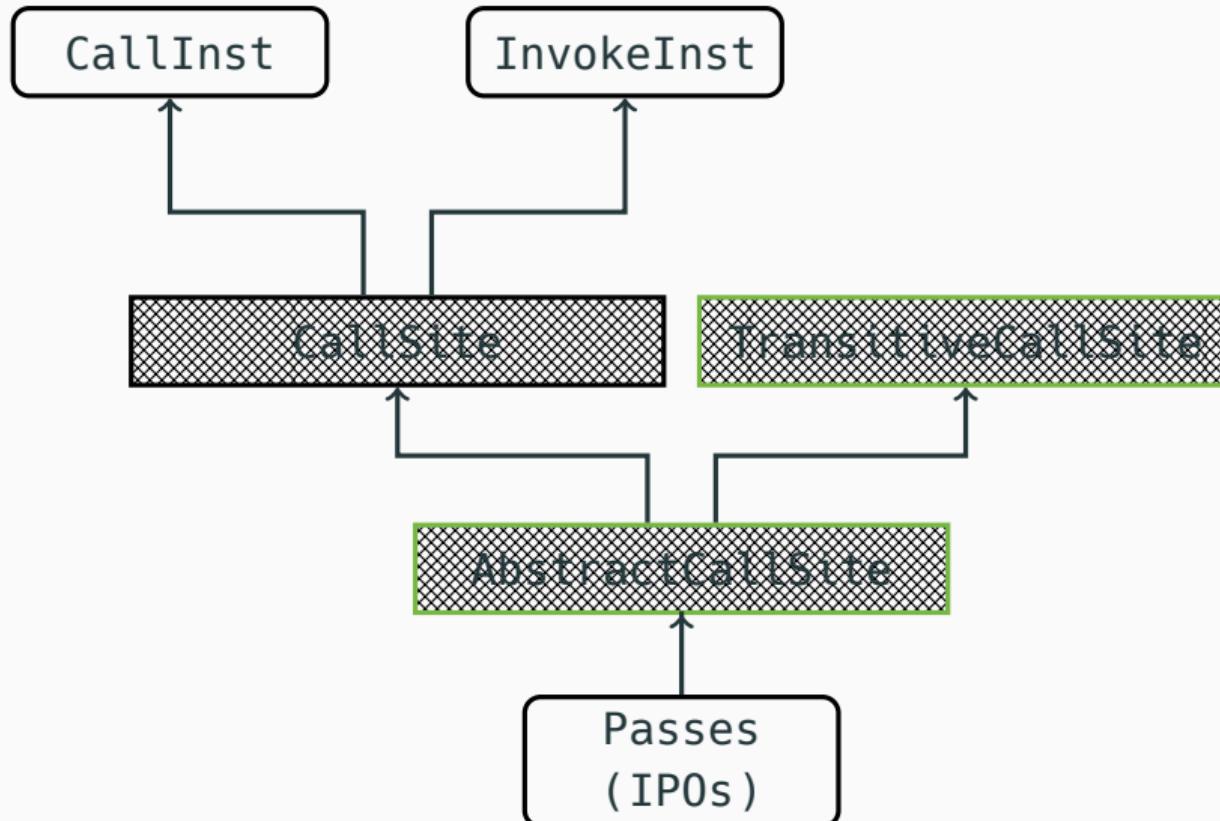
- use optimized state-machines when unavoidable
- reachability & post-dominance restrict the set of potential next parallel regions to work on
- reuse already communicated/shared values if possible
- currently, a simple state machine is generated with explicit conditionals for all known parallel regions in the module



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Call

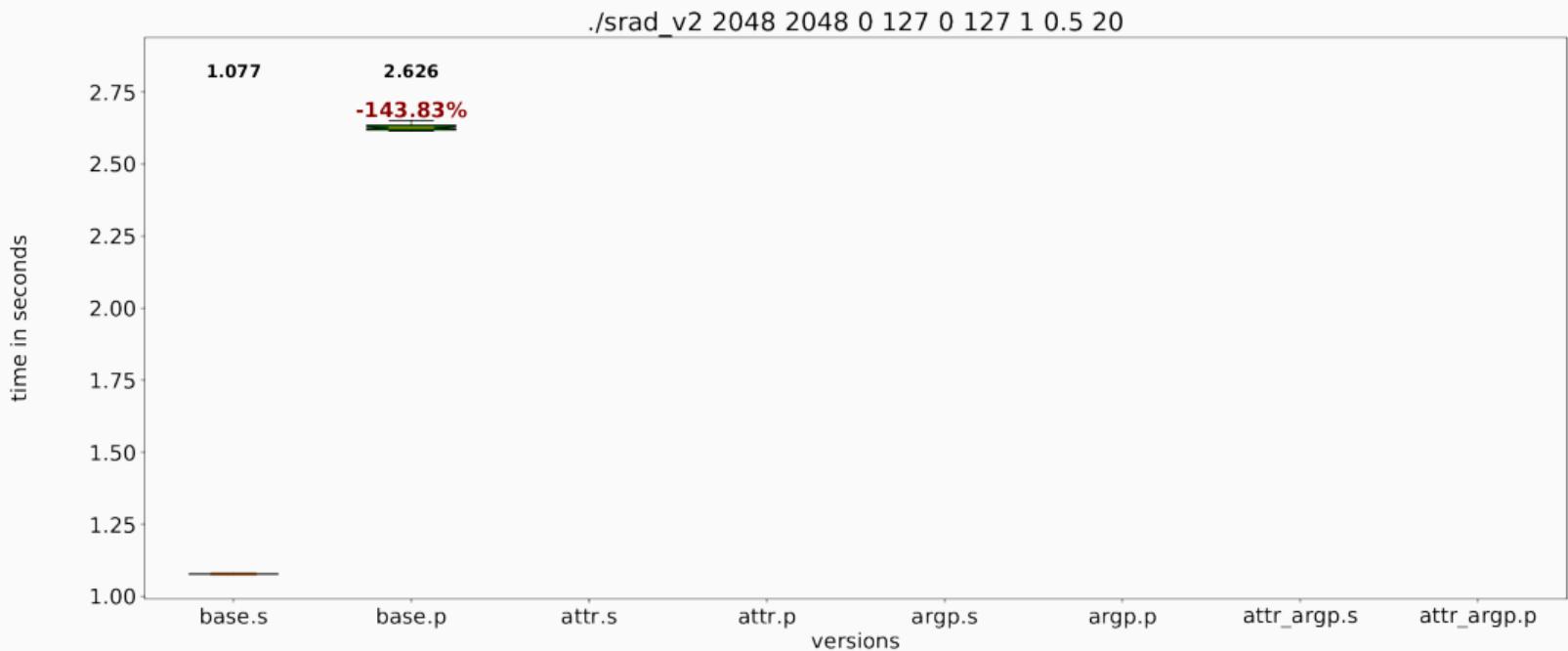
Functional changes required for  
Inter-procedural Constant Propagation:

```
for (int i = 0; i < NumArgs; i++) {  
    Value *ArgOp = ACS.getArgOperand(i);  
    if (!ArgOp) {  
        // handle non-constant  
        continue;  
    }  
    ...  
}
```

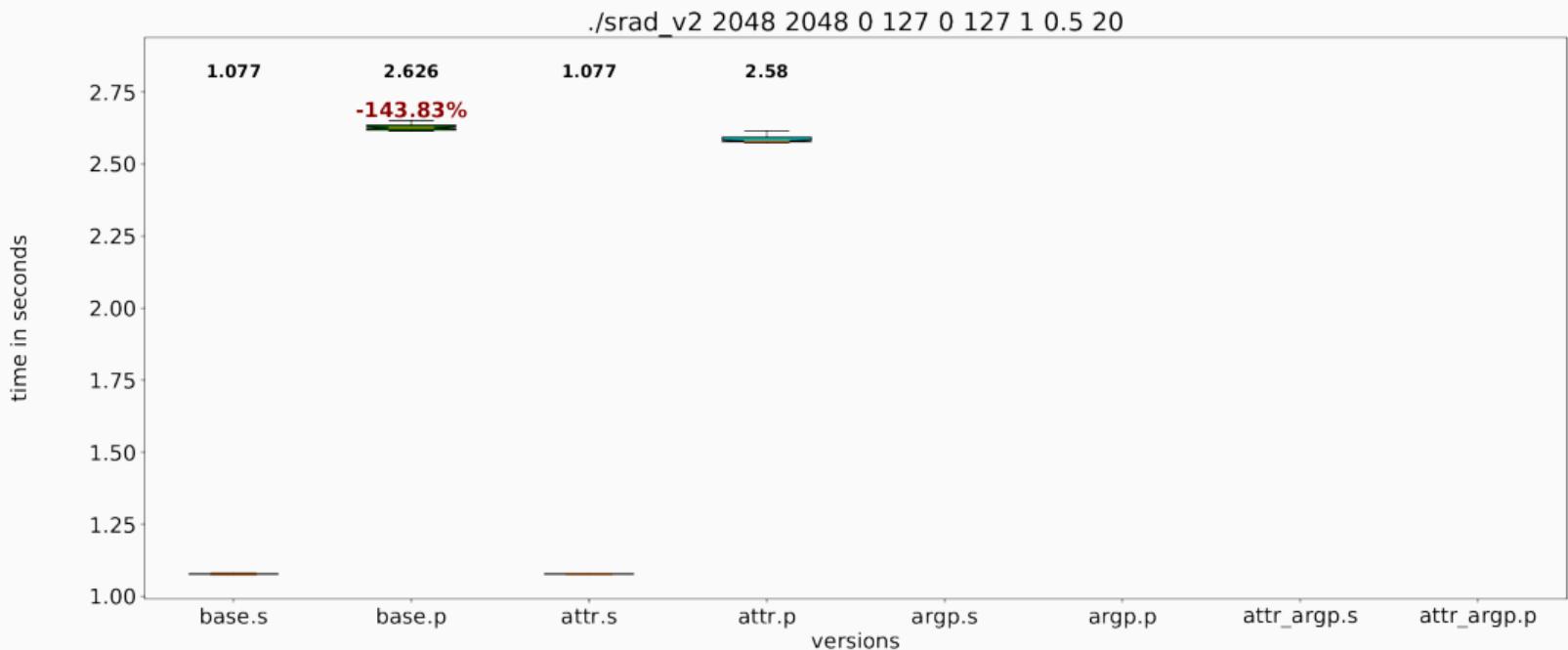
passes  
(IP0s)



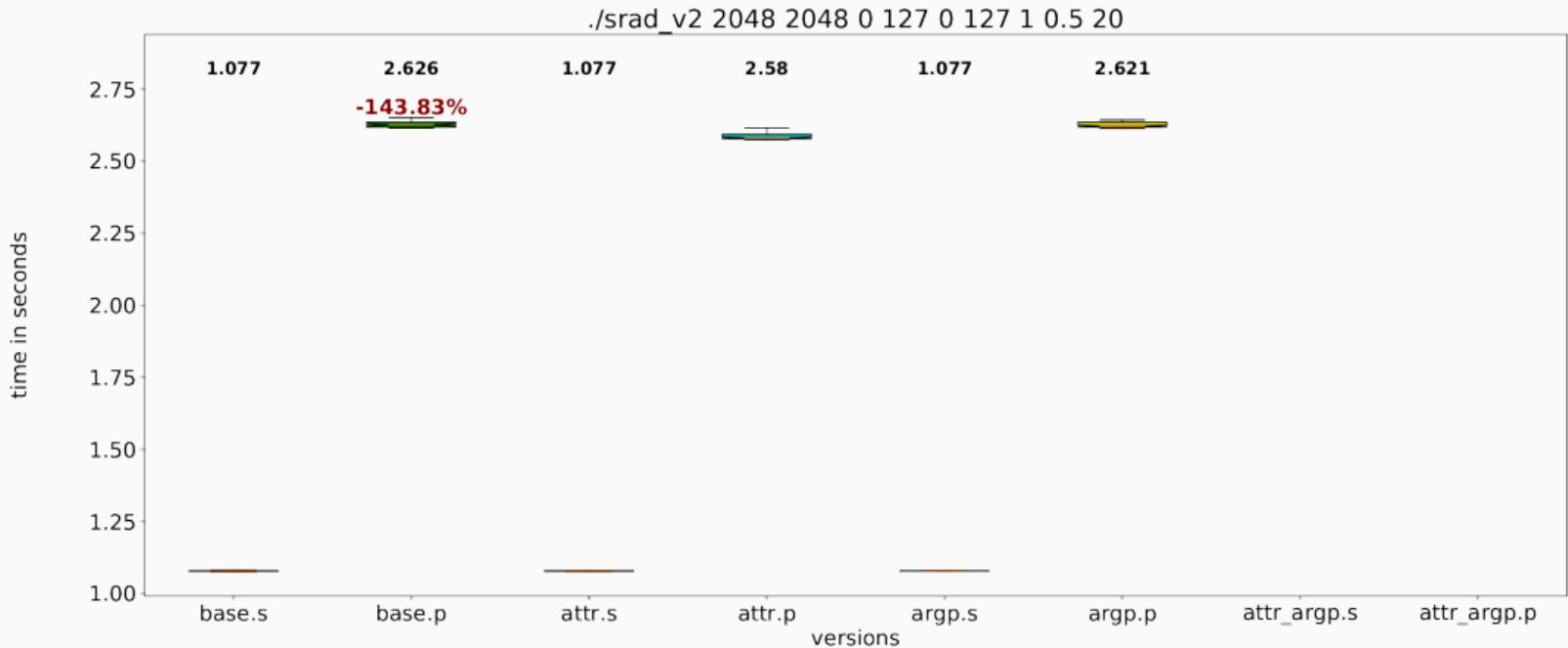
# ABSTRACT CALL SITES — PERFORMANCE RESULTS



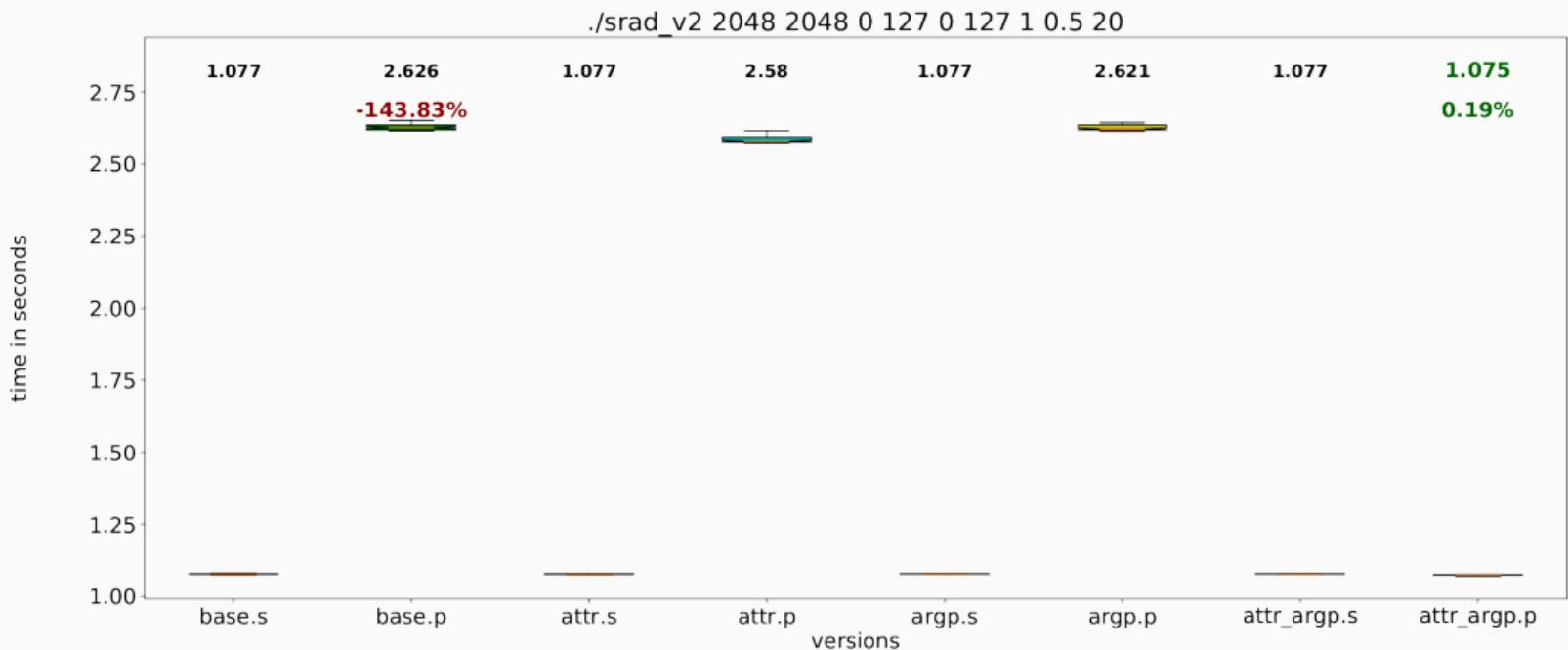
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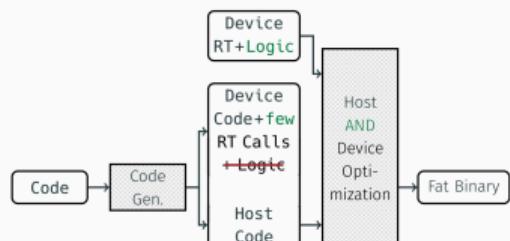
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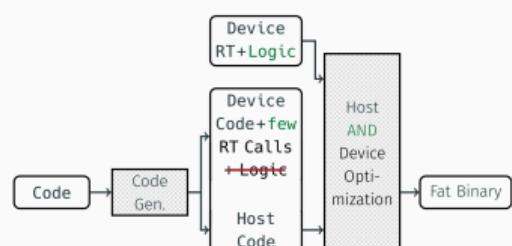
probably poor performance :(



## OPENMP OFFLOAD — OVERVIEW & DIRECTIONS

### 2. Optimize Device and Host Code Together

OpenMP    Clang    LLVM-IR    LLVM    Assembly



## OPTIMIZATION CATEGORIES

### Optimizations for sequential aspects

- May reuse existing transformations (patches up for review)  
⇒ Introduce suitable abstractions to bridge the indirection (DONE!)

### Optimizations for parallel aspects

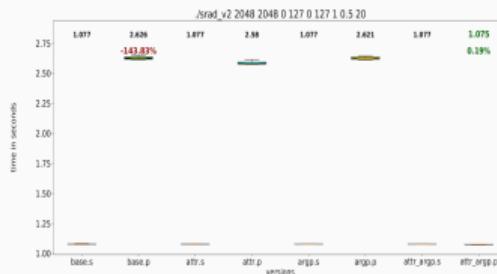
- New explicit parallelism-aware transformations (see IWOMP<sup>a</sup>)  
⇒ Introduce a unifying abstraction layer (see EuroLLVM<sup>b</sup>)

<sup>a</sup>Compiler Optimizations For OpenMP, J. Doerfert, H. Finkel, IWOMP 2018

<sup>b</sup>A Parallel IR in Real Life: Optimizing OpenMP, H. Finkel, J. Doerfert, X. Tian, G. Stelle, Euro LLVM Meeting 2018



## ABSTRACT CALL SITES — PERFORMANCE RESULTS







I: Attribute Propagation – Bidirectional Information Transfer:  
`read/write-only, restrict/noalias, ...`



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II: Variable Privatization – Limit Variable Lifetimes:

`shared(var) → firstprivate(var) → private(var)`



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III: Parallel Region Expansion – Maximize Parallel Contexts:

⇒ reduce start/stop overheads and expose barriers



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IV: Barrier Elimination – Eliminate Redundant Barriers



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comm. **vs.** operand comm. & par. compute



I: Attribute Propagation – Bidirectional Information Transfer:

**read/write-only, restrict/noalias, ...**

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comm. **vs.** operand comm. & par. compute



I: Attribute Propagation – In LLVM: Attribute Deduction (IPO!)

read/write-only, **restrict/noalias**, ...

II: Variable Privatization – In LLVM: Argument Promotion (IPO!)

**shared(var)** → **firstprivate(var)** → **private(var)**

III: Parallel Region Expansion – Maximize Parallel Contexts:

⇒ reduce start/stop overheads and expose barriers

IV: Barrier Elimination – Eliminate Redundant Barriers

V: Communication Optimization – Move Computations Around: seq. compute & result

comm. **vs.** operand comm. & par. compute



## EARLY OUTLINING

```
#pragma omp parallel for  
OpenMP Input: for (int i = 0; i < N; i++)  
              Out[i] = In[i] + In[i+N];
```

---



## EARLY OUTLINING

*#pragma omp parallel for*

OpenMP Input: **for** (**int** i = 0; i < N; i++)  
                 Out[i] = In[i] + In[i+N];

---

*// Parallel region replaced by a runtime call.*

omp\_rt\_parallel\_for(0, N, &body\_fn, &N, &In, &Out);



## EARLY OUTLINING

```
#pragma omp parallel for  
OpenMP Input: for (int i = 0; i < N; i++)  
              Out[i] = In[i] + In[i+N];
```

---

```
// Parallel region replaced by a runtime call.  
omp_rt_parallel_for(0, N, &body_fn, &N, &In, &Out);
```

```
// Parallel region outlined in the front-end (clang)!  
static void body_fn(int tid, int *N,  
                     float** In, float** Out) {  
    int lb = omp_get_lb(tid), ub = omp_get_ub(tid);  
    for (int i = lb; i < ub; i++)  
        (*Out)[i] = (*In)[i] + (*In)[i + (*N)]  
}
```



## EARLY OUTLINING

```
#pragma omp parallel for
```

```
OpenMP Input: for (int i = 0; i < N; i++)
                Out[i] = In[i] + In[i+N];
```

---

```
// Parallel region replaced by a runtime call.
```

```
omp_rt_parallel_for(0, N, &body_fn, &N, &In, &Out);
```

```
// Parallel region outlined in the front-end (clang)!
```

```
static void body_fn(int tid, int* N,
                    float** In, float** Out) {
    int lb = omp_get_lb(tid), ub = omp_get_ub(tid);
    for (int i = lb; i < ub; i++)
        (*Out)[i] = (*In)[i] + (*In)[i + (*N)]
}
```



## AN ABSTRACT PARALLEL IR

```
#pragma omp parallel for
OpenMP Input: for (int i = 0; i < N; i++)
                Out[i] = In[i] + In[i+N];
```

---

```
// Parallel region replaced by an annotated loop
for /* parallel */ (int i = 0; i < N; i++)
    body_fn(i, &N, &In, &Out);
```

```
// Parallel region outlined in the front-end (clang)!
static void body_fn(int i , int* N,
                    float** In, float** Out) {
    (*Out)[i] = (*In)[i] + (*In)[i + (*N)]
}
```



## EARLY OUTLINED + TRANSITIVE CALLS

```
#pragma omp parallel for  
OpenMP Input: for (int i = 0; i < N; i++)  
              Out[i] = In[i] + In[i+N];
```

---

```
// Parallel region replaced by a runtime call.  
omp_rt_parallel_for(0, N, &body_fn, &N, &In, &Out);  
// Model transitive call: body_fn(?, &N, &In, &Out);  
  
// Parallel region outlined in the front-end (clang)!  
static void body_fn(int tid, int* N,  
                    float** In, float** Out) {  
    int lb = omp_get_lb(tid), ub = omp_get_ub(tid);  
    for (int i = lb; i < ub; i++)  
        (*Out)[i] = (*In)[i] + (*In)[i + (*N)]  
}
```



## EARLY OUTLINED + TRANSITIVE CALLS

```
#pragma omp parallel for  
OpenMP Input: for (int i = 0; i < N; i++)  
              Out[i] = In[i] + In[i+N];
```

---

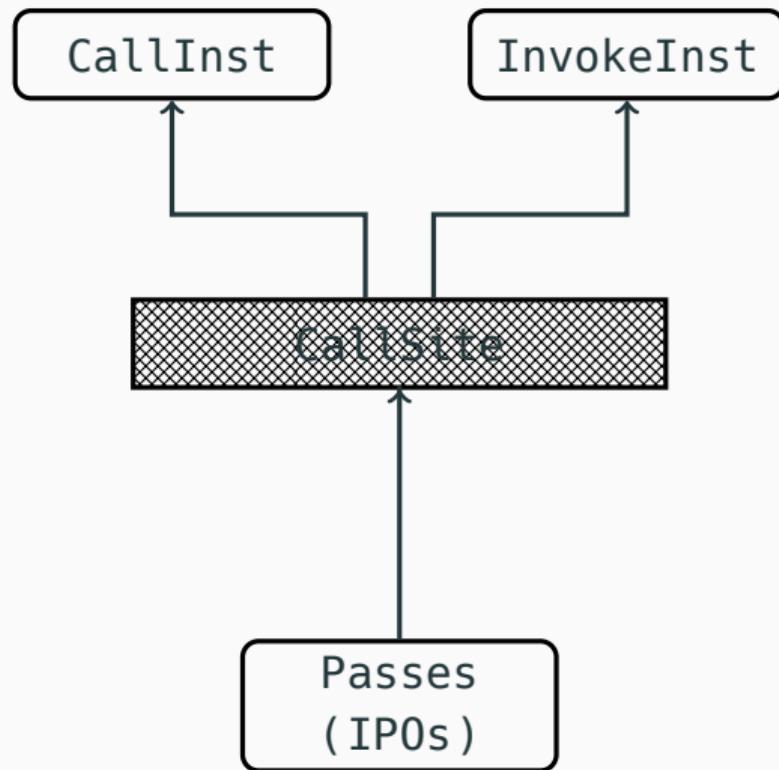
```
// Parallel region replaced by a runtime call.  
omp_rt_parallel_for(0, N, &body_fn, &N, &In, &Out);  
// Model transitive call: body_fn(?, &N, &In, &Out);
```

```
// Parallel region outlined in the front-end (clang)!  
static void body_fn(int tid, int* N,  
                    float** In, float** Out) {  
    int lb = 0;  
    for (int i = lb; i < N[0]; i++)  
        (*Out)[i] = (*In)[i] + (*In)[i+N[0]] * tid;  
}
```

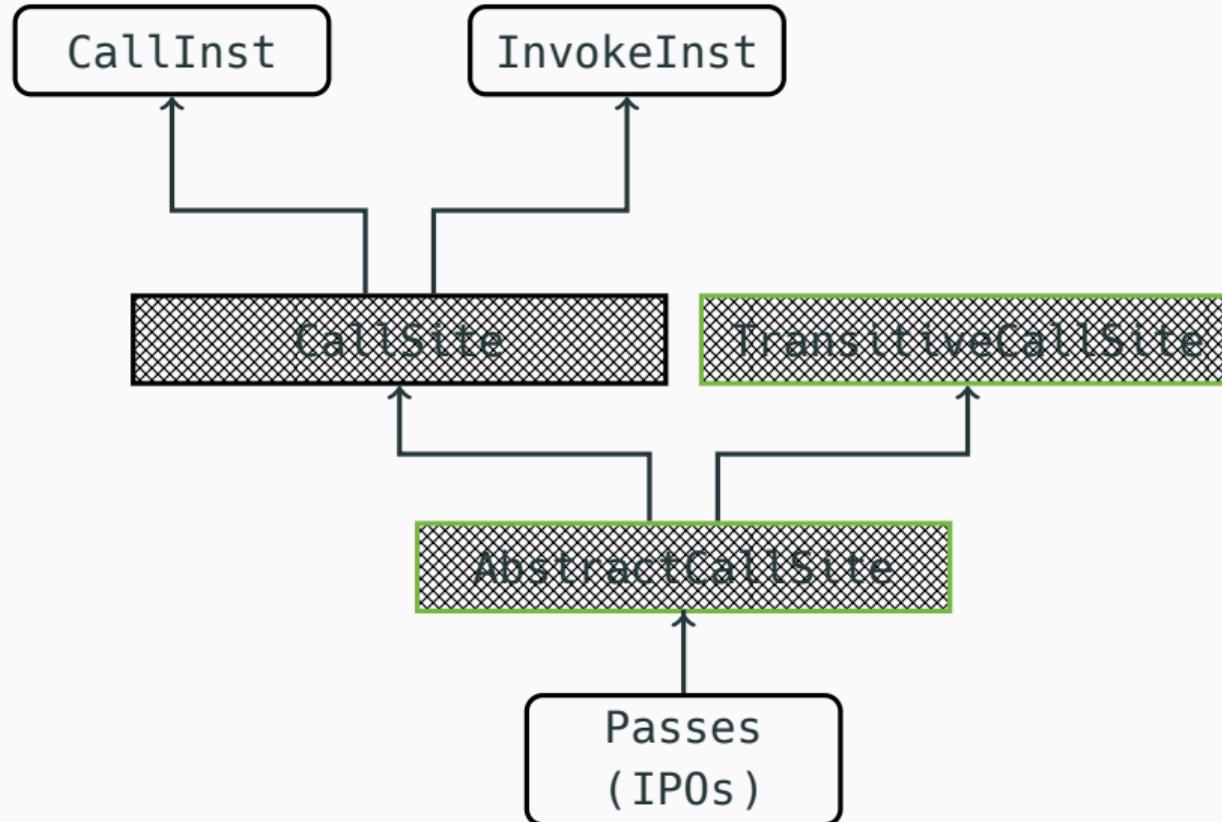
- + valid and executable IR
- integration cost per IPO
- + no unintended interactions



# IPO IN LLVM



# TRANSITIVE CALL SITES IN LLVM



## TRANSITIVE CALL SITES IN LLVM

Call

Functional changes required for  
Inter-procedural Constant Propagation:

```
for (int i = 0; i < NumArgs; i++) {  
    Value *ArgOp = ACS.getArgOperand(i);  
    if (!ArgOp) {  
        // handle non-constant  
        continue;  
    }  
    ...  
}
```

(IP0s)



## EVALUATED VERSION

Version	Description	Opt.
<i>base</i>	plain “-O3”, thus no parallel optimizations	
<i>attr</i>	attribute propagation through attr. deduction (IPO)	I
<i>argp</i>	variable privatization through arg. promotion (IPO)	II
<i>n/a</i>	constant propagation (IPO)	





## SOME CONTEXT



## SOME CONTEXT

### Examples

Examples are given in a C-like language with OpenMP annotations.

### Transformations

Our transformations work on the LLVM intermediate representation (LLVM-IR), thus take and produce LLVM-IR.

### OpenMP Runtime Library

We experience OpenMP annotations as OpenMP runtime library calls and the situation is most often more complicated than presented here.



## EVALUATION ENVIRONMENT

- Run with 1 Thread<sup>2</sup>
- Median and variance of 51 runs is shown
- Rodiana 3.1 benchmarks and LULESH v1.0 (OpenMP)
- Only time in parallel constructs was measured

---

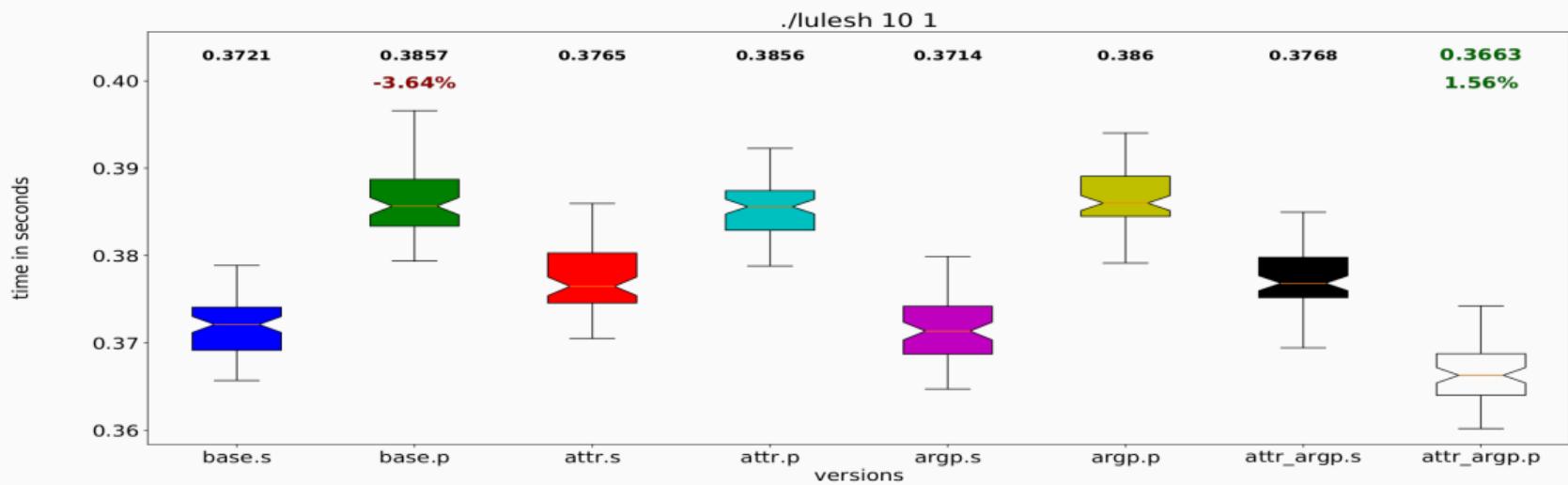
<sup>2</sup>Intel(R) Core(TM) i7-4800MQ CPU @ 2.70GHz



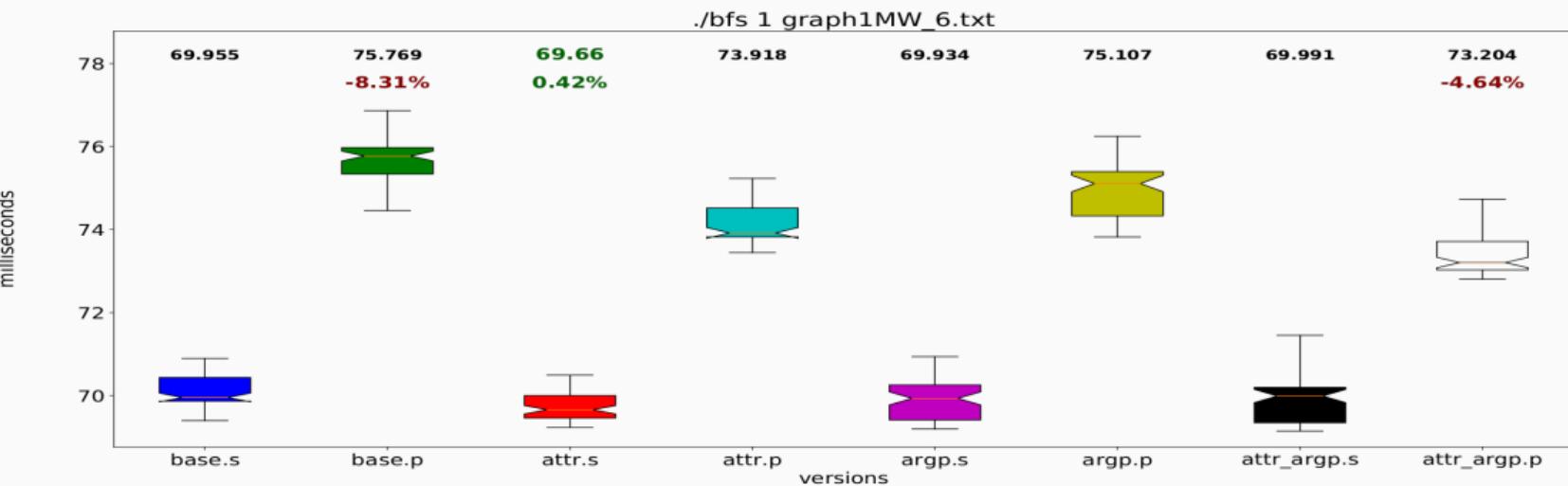
# PERFORMANCE RESULTS



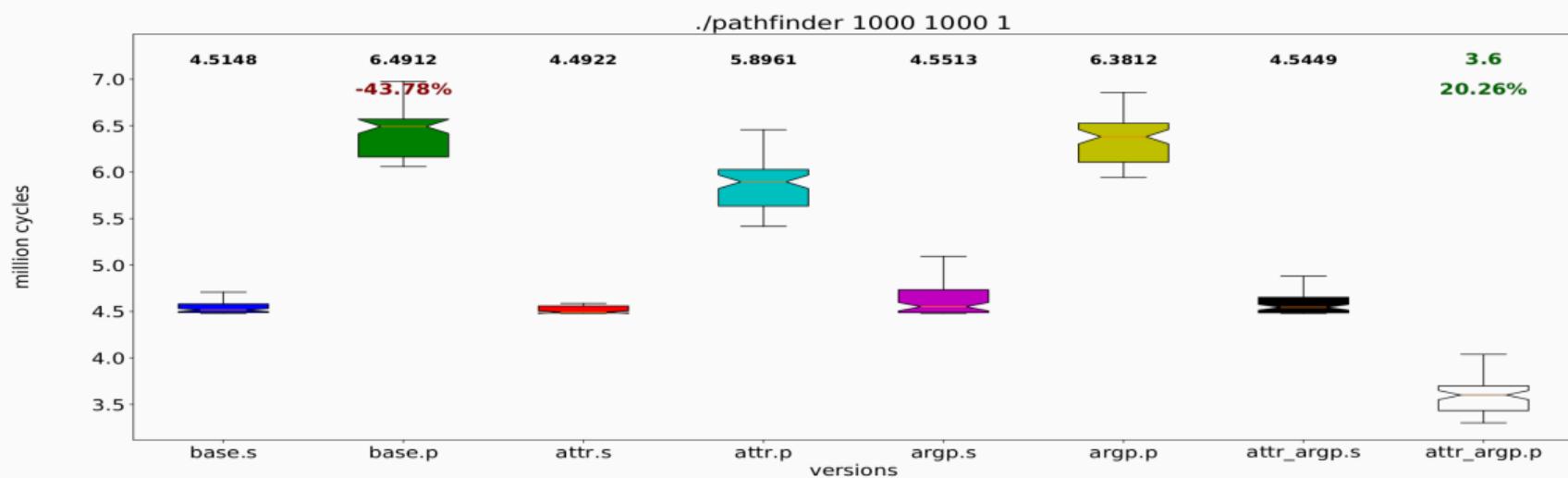
# PERFORMANCE RESULTS



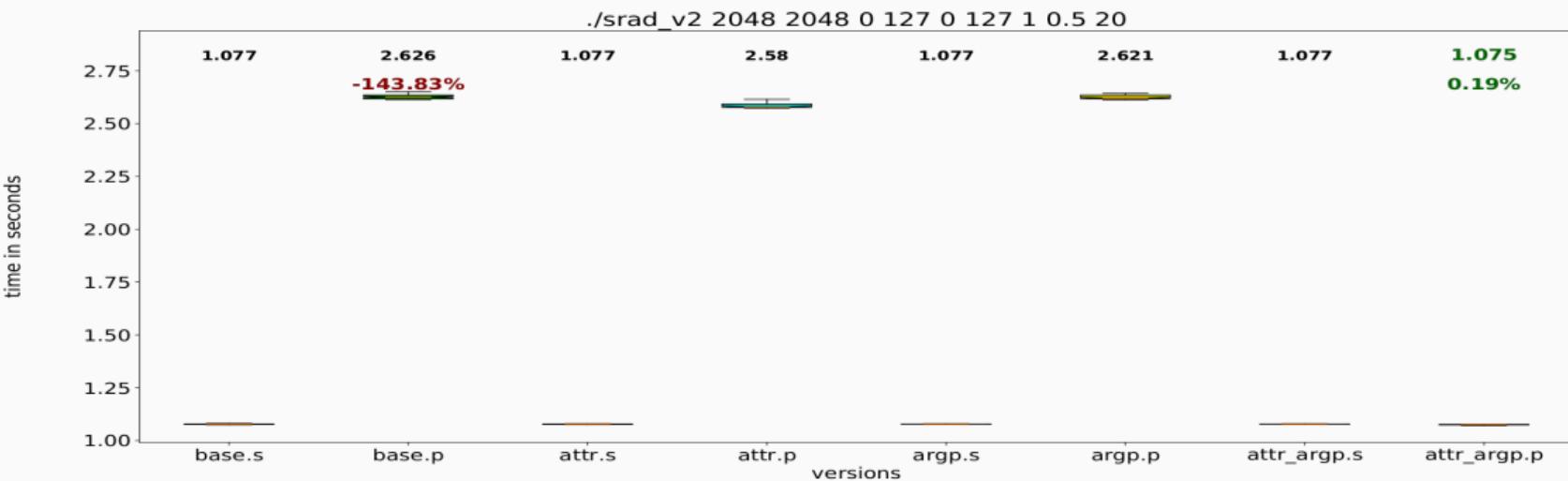
# PERFORMANCE RESULTS



# PERFORMANCE RESULTS



# PERFORMANCE RESULTS



## ACTION ITEM I

---

\*

†



## ACTION ITEM I

- 1) Run *your* OpenMP code sequentially<sup>\*</sup>,  
with and without OpenMP.

---

<sup>\*</sup>export OMP\_NUM\_THREADS=1  
†



## ACTION ITEM I

- 1) Run *your* OpenMP code sequentially<sup>\*</sup>,  
with and without OpenMP.
  
- 2) Email me<sup>†</sup> the results!

---

<sup>\*</sup>export OMP\_NUM\_THREADS=1

<sup>†</sup>jdoerfert@anl.gov



## ACTION ITEM II

---

★



## ACTION ITEM II

- 1) Always<sup>\*</sup> use `default(none)` and `firstprivate(...)`

---

<sup>\*</sup>For scalars/pointers if you do not have explicit synchronization.



## ACTION ITEM II

- 1) Always<sup>\*</sup> use `default(none)` and  
`firstprivate(...)`
- 2) Revisit ACTION ITEM I

---

<sup>\*</sup>For scalars/pointers if you do not have explicit synchronization.



## ACTION ITEM II

NO need to “share” the variable A!

```
#pragma omp parallel for shared(A)
for (int i = 0; i < N; i++)
    A[i] = i;
```

---

\* For scalars/pointers if you do not have explicit synchronization.



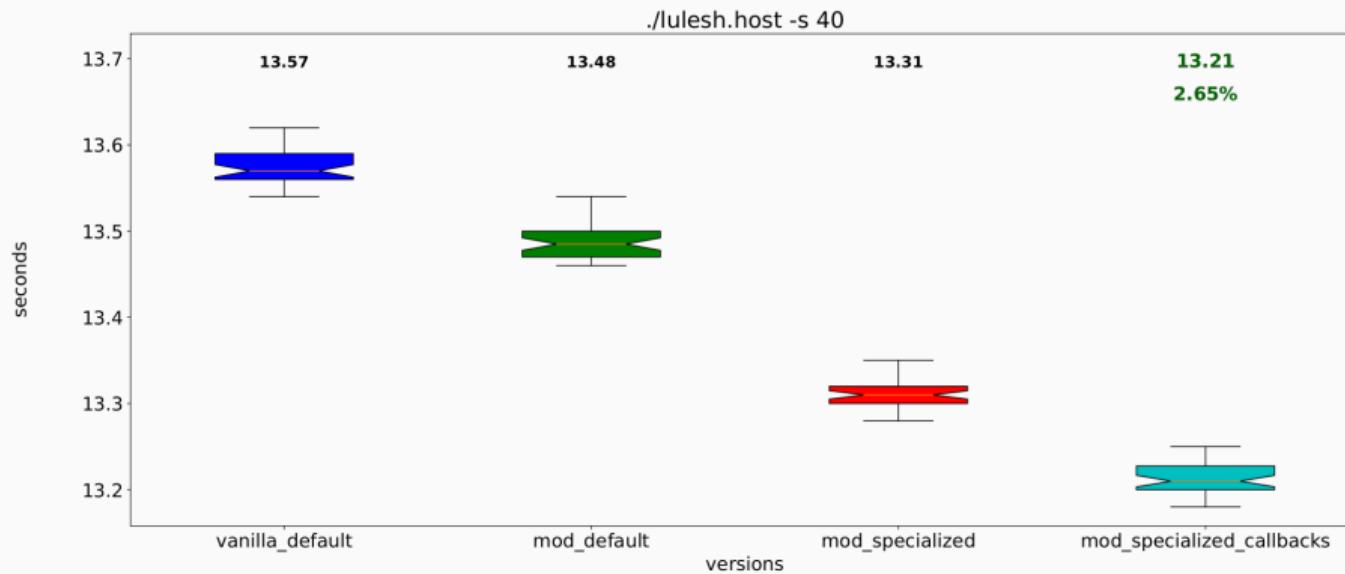
## CONSTANT PROPAGATION EXAMPLE

```
double gamma[4][8];
gamma[0][0] = 1;
// ... and so on till ...
gamma[3][7] = -1;
```

```
Kokkos::parallel_for(
    "CalcFBHourglassForceForElems A",
    numElem, KOKKOS_LAMBDA(const int &i2) {
        // Use gamma[0][0] ... gamme[3][7]
    }
```



# CONSTANT PROPAGATION PERFORMANCE



## OPTIMIZATION I: ATTRIBUTE PROPAGATION

*#pragma omp parallel for firstprivate(...)*

OpenMP Input: **for** (**int** i = 0; i < N; i++)  
    Out[i] = In[i] + In[i+N];

---



## OPTIMIZATION I: ATTRIBUTE PROPAGATION

```
#pragma omp parallel for firstprivate(...)
```

OpenMP Input: **for** (**int** i = 0; i < N; i++)  
    Out[i] = In[i] + In[i+N];

---

```
// Parallel region replaced by a runtime call.  
omp_rt_parallel_for(0, N, &body_fn, N, In, Out);
```

```
// Parallel region outlined in the front-end (clang)!  
void body_fn(int i, int N,  
              float* In,  
              float* Out) {  
    Out[i] = In[i] + In[i + N];  
}
```



## OPTIMIZATION I: ATTRIBUTE PROPAGATION

*#pragma omp parallel for firstprivate(...)*

OpenMP Input: **for** (**int** i = 0; i < N; i++)  
    Out[i] = In[i] + In[i+N];

---

*// Parallel region replaced by a runtime call.*  
**omp\_rt\_parallel\_for(0, N, &body\_fn, N, In, Out);**

*// Parallel region outlined in the front-end (clang)!*  
**void** body\_fn(**int** i, **int** N,  
              **float**\* /\* read-only & no-escape \*/ In,  
              **float**\* /\* write-only & no-escape \*/ Out) {  
    Out[i] = In[i] + In[i + N];  
}



## OPTIMIZATION I: ATTRIBUTE PROPAGATION

*#pragma omp parallel for firstprivate(...)*

OpenMP Input: **for** (**int** i = 0; i < N; i++)  
    Out[i] = In[i] + In[i+N];

---

*// Parallel region replaced by a runtime call.*

```
omp_rt_parallel_for(0, N, &body_fn, N,  
    /* ro & no-esc */ In, /* wo & no-esc */ Out);
```

*// Parallel region outlined in the front-end (clang)!*

```
void body_fn(int i, int N,  
            float* /* read-only & no-escape */ In,  
            float* /* write-only & no-escape */ Out) {  
    Out[i] = In[i] + In[i + N];  
}
```



```
int foo() {
    int a = 0;

#pragma omp parallel
{
#pragma omp critical
{ a += 1; }
bar();
#pragma omp critical
{ a *= 2; }
}
return a;
}
```

```
int foo() {  
    int a = 0;  
  
#pragma omp parallel  
{  
#pragma omp critical  
{ a += 1; }  
bar();  
#pragma omp critical  
{ a *= 2; }  
}  
return a;  
}
```

```
int foo() {  
    int a = 0;  
    int *restrict p = &a;  
    omp_rt_parallel_for(pwork, p);  
    return a;  
}  
void pwork(int tid, int *p) {  
    if (omp_critical(tid)) {  
        *p = *p + 1;  
        omp_critical_end(tid);  
    }  
    bar();  
    if (omp_critical(tid)) {  
        *p = *p * 2;  
        omp_critical_end(tid);  
    }  
}
```

```
void pwork(int tid,
           int *restrict p) {
    if (omp_critical(tid)) {
        omp_critical_end(tid);
    }
    bar();
    if (omp_critical(tid)) {
        *p = 2 * (*p + 1);
        omp_critical_end(tid);
    }
}
```

```
int foo() {
    int a = 0;
    int *restrict p = &a;
    omp_rt_parallel_for(pwork, p);
    return a;
}

void pwork(int tid, int *p) {
    if (omp_critical(tid)) {
        *p = *p + 1;
        omp_critical_end(tid);
    }
    bar();
    if (omp_critical(tid)) {
        *p = *p * 2;
        omp_critical_end(tid);
    }
}
```

```
void pwork(int tid,
           int *restrict p) {
    if (omp_critical(tid)) {
        *p = *p + 1;
        omp_critical_end(tid);
    }
    bar()[p]; // May "use" p.
    if (omp_critical(tid)) {
        *p = *p * 2;
        omp_critical_end(tid);
    }
}
```

```
int foo() {
    int a = 0;
    int *restrict p = &a;
    omp_rt_parallel_for(pwork, p);
    return a;
}

void pwork(int tid, int *p) {
    if (omp_critical(tid)) {
        *p = *p + 1;
        omp_critical_end(tid);
    }
    bar();
    if (omp_critical(tid)) {
        *p = *p * 2;
        omp_critical_end(tid);
    }
}
```

## OPTIMIZATION II: VARIABLE PRIVATIZATION

*#pragma omp parallel for shared(...)*

OpenMP Input: **for** (**int** i = 0; i < N; i++)  
    Out[i] = In[i] + In[i+N];

---



## OPTIMIZATION II: VARIABLE PRIVATIZATION

*#pragma omp parallel for shared(...)*

OpenMP Input: **for** (**int** i = 0; i < N; i++)  
    Out[i] = In[i] + In[i+N];

---

*// Parallel region replaced by a runtime call.*

```
omp_rt_parallel_for(0, N, &body_fn, &N, &In, &Out);
```

*// Parallel region outlined in the front-end (clang)!*

```
void body_fn(int i, int* N,  
             float** In,  
             float** Out) {  
    (*Out)[i] = (*In)[i] + (*In)[i + (*N)];  
}
```



## OPTIMIZATION II: VARIABLE PRIVATIZATION

*#pragma omp parallel for shared(...)*

OpenMP Input: **for** (**int** i = 0; i < N; i++)  
    Out[i] = In[i] + In[i+N];

---

*// Parallel region replaced by a runtime call.*

```
omp_rt_parallel_for(0, N, &body_fn, &N, &In, &Out);
```

*// Parallel region outlined in the front-end (clang)!*

```
void body_fn(int i, int* /* ro & ne */ N,
             float** /* ro & ne */ In,
             float** /* ro & ne */ Out) {
    (*Out)[i] = (*In)[i] + (*In)[i + (*N)];
}
```



## OPTIMIZATION II: VARIABLE PRIVATIZATION

```
#pragma omp parallel for firstprivate(...)
```

OpenMP Input: **for** (**int** i = 0; i < N; i++)  
    Out[i] = In[i] + In[i+N];

---

```
// Parallel region replaced by a runtime call.  
omp_rt_parallel_for(0, N, &body_fn, N, In, Out);
```

```
// Parallel region outlined in the front-end (clang)!  
void body_fn(int i, int N,  
              float* In,  
              float* Out) {  
    Out[i] = In[i] + In[i + N];  
}
```



## OPTIMIZATION III: PARALLEL REGION EXPANSION



## OPTIMIZATION III: PARALLEL REGION EXPANSION

```
void copy(float* dst, float* src, int N) {
    #pragma omp parallel for
    for(int i = 0; i < N; i++) {
        dst[i] = src[i];
    } // implicit barrier!
}

void compute_step_factor(int nelr, float* vars,
                        float* areas, float* sf) {
    #pragma omp parallel for
    for (int blk = 0; blk < nelr / block_length; ++blk) {
        ...
    } // implicit barrier!
}
```



## OPTIMIZATION III: PARALLEL REGION EXPANSION

```
for (int i = 0; i < iterations; i++) {  
    copy(old_vars, vars, nelr * NVAR);  
  
    compute_step_factor(nelr, vars, areas, sf);  
  
    for (int j = 0; j < RK; j++) {  
        compute_flux(nelr, ese, normals, vars, fluxes, ff_vars,  
                     ff_m_x, ff_m_y, ff_m_z, ff_dnergy);  
  
        time_step(j, nelr, old_vars, vars, sf, fluxes);  
    }  
}
```



## OPTIMIZATION III: PARALLEL REGION EXPANSION

```
for (int i = 0; i < iterations; i++) {  
    #pragma omp parallel for          // copy  
    for (...) {  
        /* write old_vars, read vars */  
    } // implicit barrier!  
    compute_step_factor(nelr, vars, areas, sf);  
  
    for (int j = 0; j < RK; j++) {  
        compute_flux(nelr, ese, normals, vars, fluxes, ff_vars,  
                     ff_m_x, ff_m_y, ff_m_z, ff_dnergy);  
  
        time_step(j, nelr, old_vars, vars, sf, fluxes);
```



## OPTIMIZATION III: PARALLEL REGION EXPANSION

```
for (int i = 0; i < iterations; i++) {  
    #pragma omp parallel for           // copy  
    for (...) {  
        /* write old_vars, read vars */  
    } // implicit barrier!  
    #pragma omp parallel for           // compute_step_factor  
    for (...) {  
        /* write sf, read vars & area */  
    } // implicit barrier!  
    for (int j = 0; j < RK; j++) {  
        #pragma omp parallel for         // compute_flux  
        for (...) {  
            /* write fluxes, read vars & ... */  
        } // implicit barrier!  
    } // implicit barrier!  
}
```



## OPTIMIZATION III: PARALLEL REGION EXPANSION

```
#pragma omp parallel
for (int i = 0; i < iterations; i++) {
    #pragma omp for                      // copy
    for (...) {
        /* write old_vars, read vars */
    } // explicit barrier in LLVM-IR!
    #pragma omp for                      // compute_step_factor
    for (...) {
        /* write sf, read vars & area */
    } // explicit barrier in LLVM-IR!
    for (int j = 0; j < RK; j++) {
        #pragma omp for                  // compute_flux
        for (...) {
            /* write fluxes, read vars & ... */
        } // explicit barrier in LLVM-IR!
    ...
}
```



## OPTIMIZATION IV: BARRIER ELIMINATION

```
#pragma omp parallel
for (int i = 0; i < iterations; i++) {
    #pragma omp for           // copy
    for (...) {
        /* write old_vars, read vars */
    } // explicit barrier in LLVM-IR!
    #pragma omp for           // compute_step_factor
    for (...) {
        /* write sf, read vars & area */
    } // explicit barrier in LLVM-IR!
    for (int j = 0; j < RK; j++) {
        #pragma omp for           // compute_flux
        for (...) {
            /* write fluxes, read vars & ... */
        } // explicit barrier in LLVM-IR!
    ...
}
```



## OPTIMIZATION IV: BARRIER ELIMINATION

```
#pragma omp parallel
for (int i = 0; i < iterations; i++) {
    #pragma omp for                      // copy
    for (...) {
        /* write old_vars, read vars */
    } // explicit barrier in LLVM-IR!
    #pragma omp for                      // compute_step_factor
    for (...) {
        /* write sf, read vars & area */
    } // explicit barrier in LLVM-IR!
    for (int j = 0; j < RK; j++) {
        #pragma omp for                  // compute_flux
        for (...) {
            /* write fluxes, read vars & ... */
        } // explicit barrier in LLVM-IR!
    ...
}
```



## OPTIMIZATION IV: BARRIER ELIMINATION

```
#pragma omp parallel
for (int i = 0; i < iterations; i++) {
    #pragma omp for nowait          // copy
    for (...) {
        /* write old_vars, read vars */
    }
    #pragma omp for nowait          // compute_step_factor
    for (...) {
        /* write sf, read vars & area */
    }
    for (int j = 0; j < RK; j++) {
        #pragma omp for              // compute_flux
        for (...) {
            /* write fluxes, read vars & ... */
        }
    } // explicit barrier in LLVM-IR!
    ...
}
```



## OPTIMIZATION V: COMMUNICATION OPTIMIZATION



## OPTIMIZATION V: COMMUNICATION OPTIMIZATION

```
void f(int *X, int *restrict Y) {
    int L = *X;           // immovable
    int N = 512;          // movable

    int A = N + L;        // movable
#pragma omp parallel for
    firstprivate(X, Y, N, L, A)
    for (int i = 0; i < N; i++) {
        int K = *Y;        // movable
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    }
}
```



# OPTIMIZATION V: COMMUNICATION OPTIMIZATION

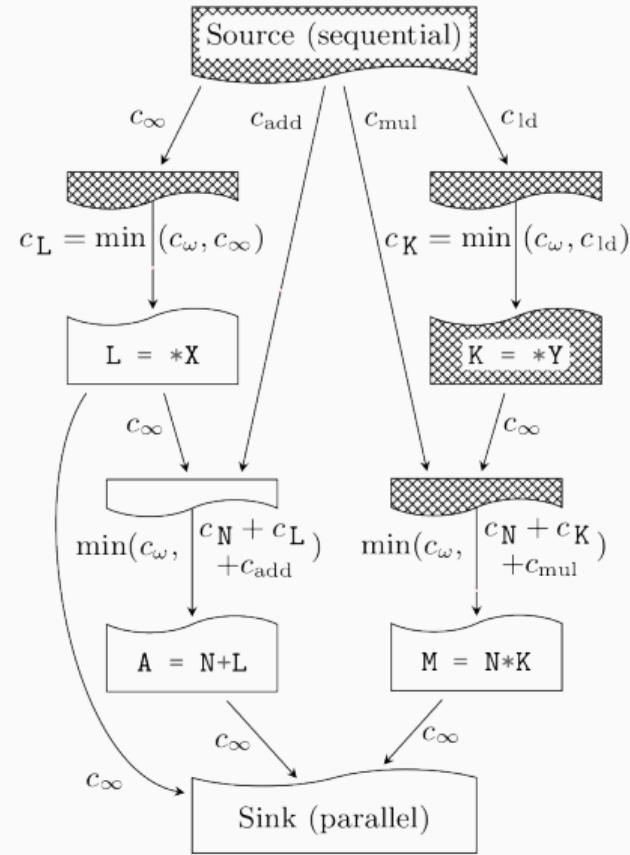
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$$\begin{array}{llll}
 c_\infty = \infty & c_{\text{add}} = 5 & c_\omega = 15 & c_{\text{mul}} = 10 \\
 c_{\text{ld}} = 20 & c_N = c_{\text{cst}} & c_{\text{cst}} = 0 & \text{cut}
 \end{array}$$

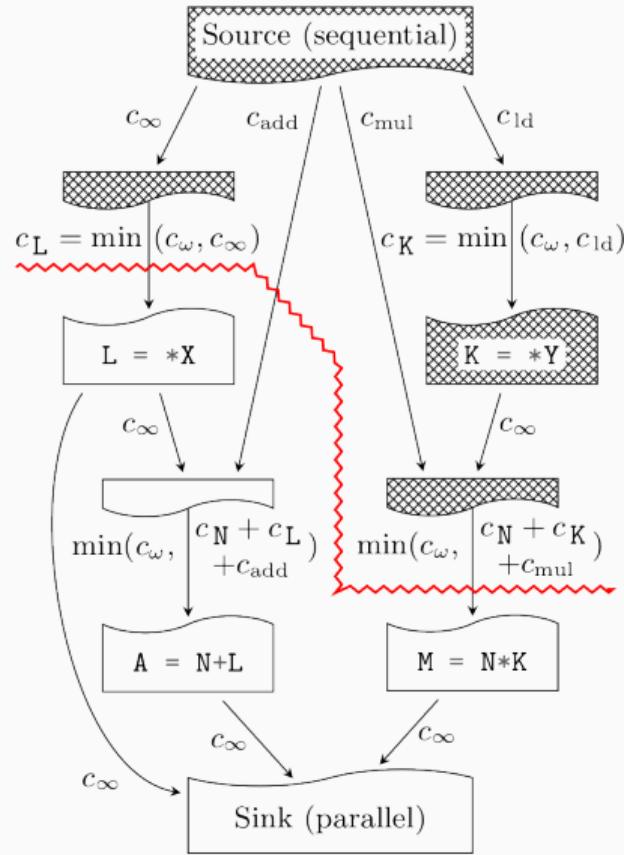


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        int K = *Y;         // movable
        int M = N * K;      // movable
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    }
}
```

```
void g(int *X, int *restrict Y) {
    int L = *X;           // immovable
    int K = *Y;           //  $c_{ld} > c_\omega$ 
    int M = 512 * K;      //  $c_{mul} + c_K > c_\omega$ 
#pragma omp parallel
    firstprivate(X, M, L)
{
    int A = 512 + L;     //  $c_{add} < c_\omega$ 
#pragma omp for
    firstprivate(X, M, A, L)
    for (int i = 0; i < 512; i++) {
        X[i] = M+A*L*i; // immovable
    }
}
```



# EARLY OUTLINING: SEQUENTIAL OPTIMIZATION PROBLEMS

Information Transfer

Value Transfer



## EARLY OUTLINING: SEQUENTIAL OPTIMIZATION PROBLEMS

NO Information Transfer:  
*outlined function  $\Leftrightarrow$  runtime library call site*

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Declaration	Value Transfer OpenMP Clause	Communication Type
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T var;	<b>shared(var)</b>	
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T var;	<i>firstprivate(var)</i>	var of type T
T var;	<i>private(var)</i>	<i>none</i>





# TARGET REGION — THE INTERFACE

```
void kernel(...) {  
  
    init:  
        char ThreadKind = __kmpc_target_region_kernel_init(...);  
        if (ThreadKind == -1) { // actual worker thread  
            if (!UsedLibraryStateMachine)  
                user_code_state_machine();  
            goto exit;  
        } else if (ThreadKind == 0) { // surplus worker thread  
            goto exit;  
        } else { // team master thread  
            goto user_code;  
        }  
  
    user_code:  
        // User defined kernel code, parallel regions are replaced by  
        // by __kmpc_target_region_kernel_parallel(...) calls.  
  
        // Fallthrough to de-initialization  
   _deinit:  
        __kmpc_target_region_kernel_deinit(...);  
  
    exit:  
        /* exit the kernel */  
}
```



## TARGET REGION — THE INTERFACE

```
// Initialization
int8_t __kmpc_target_region_kernel_init(ident_t *Ident,
                                         bool UseSPMDMode,
                                         bool RequiresOMPRuntime,
                                         bool UseStateMachine,
                                         bool RequiresDataSharing);

// De-Initialization
void __kmpc_target_region_kernel_deinit(ident_t *Ident,
                                         bool UseSPMDMode,
                                         bool RequiredOMPRuntime);

// Parallel execution
typedef void (*ParallelWorkFnTy)(void * /* SharedValues */,
                                    void * /* PrivateValues */)

CALLBACK(ParallelWorkFnTy, SharedValues, PrivateValues)
void __kmpc_target_region_kernel_parallel(ident_t *Ident,
                                         bool UseSPMDMode, bool RequiredOMPRuntime,
                                         ParallelWorkFnTy ParallelWorkFn, void *SharedValues,
                                         uint16_t SharedValuesBytes, void *PrivateValues,
                                         uint16_t PrivateValuesBytes, bool SharedMemPointers);
```



## TARGET REGION — THE IMPLEMENTATION

- (almost) the same as with the current NVPTX backend, except for shared/firstprivate variables
- implemented in Cuda as part of the library, not generated into the user code/module/TU by Clang
- the boolean flags are commonly constant, after inlining all target region abstractions is gone



## ACTION ITEM III

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†



## ACTION ITEM III

- 1) Review your OpenMP target code.

---

†



## ACTION ITEM III

- 1) Review your OpenMP target code.
- 2) Email me<sup>†</sup> if you use the “bad” pattern!

---

<sup>†</sup>[jdoerfert@anl.gov](mailto:jdoerfert@anl.gov)



## CURRENT WORK — REVIEWS, EVALUATION, FEATURES, HARDENING

- started the review process



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